Hash Functions, Message Authentication Codes

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Security Services
- Confidentiality: Symmetric encryption solves
- Integrity
- Authentication
- Non-repudiation
- Access control
- Availability

Integrity in Networking
- Sender computes a CRC for the message
- Sender appends the CRC code to the message and sends them to the receiver
- The receiver computes the CRC of the message.
  - If the CRC appended to the message is equal to the computed one, the message is unchanged with a high probability.
  - If the CRCs do not match, the message is changed during transmission.

CRC Checksum in Networking

Sender

Receiver

INTERNET

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<thead>
<tr>
<th>M</th>
<th>CRC</th>
<th>Chk Sum</th>
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Cryptographic Hash Functions

- Maps an arbitrary length input to a fixed-size output.
  - If \( m \) is message, \( H \) is the hash function, \( H(m) \) is the output of hash function, also called message digest.
- Desirable features:
  - One-way: There should be no easy way to guess \( m \) from \( H(m) \)
  - Pseudorandom: If \( m \) and \( m' \) are two close values, \( H(m) \) and \( H(m') \) should not be close each other.
  - Collision resistant: It should be hard to find two inputs that hash to the same output
    - It should be hard to find two inputs \( a \) and \( b \) such that \( H(a) = H(b) \)

Example Operation of Hash Functions

Birthday Paradox

- Birthday Problem ("paradox"): When \( \sqrt{N} \) or more are chosen randomly from a domain of \( N \), there is a significant chance of collision.
- Probability of \( n \) persons having different birthdays:
  \[
p(n) = 1 \times (1 - \frac{1}{365}) \times (1 - \frac{2}{365}) \times \ldots \times (1 - \frac{n-1}{365})
\]
Collision Resistance

- If a hash function produces \( N \) bits of output, an attacker should not easily find a collision by performing less than (on average) \( 2^{N/2} \) hash operations.
  - If there is an easier method than this brute force attack, it is typically considered a flaw in the hash function.
  - Therefore, hash output size \( \geq 128 \) bits is desirable.
- But why “collision resistance”?
  - A chosen plaintext attack: Trudy is Alice’s secretary. Generates two opposite messages.

Internals of a Hash Function

- A fixed-size “compression function”.
  - Each iteration mixes an input block with the previous output.
  - Design:
    - Lots of operations (rotations, \( \oplus \), \( \land \), \( \lor \), +,...) fast in s/w.
    - More of them are added if a weakness is found.

Some Popular Hash Algorithms

- MD5 (Rivest)
  - 128-bit output
  - Most popular
- SHA-1 (NIST-NSA)
  - US gov’t standard
  - 160-bit output
- RIPEMD-160
  - Euro. RIPE project.
  - 160-bit output

Message Authentication Codes (MAC)

- A simple message integrity checking method:
  - Compute \( H(m) \) and send \((m, H(m))\)
  - The receiver computes \( H(m) \) and compares with the received \( H(m) \) value.
- What happens if an attacker changes both \( m \) and \( H(m) \) value and sends \((m', H(m'))\) to receiver?
  - A secret key system can be used to generate a cryptographic checksum known as a message authentication code (MAC).
    - It is also referred as MIC (Message Integrity Code).
MACs

- Let $\text{MAC}_K(m)$ be a message authentication code for $m$ produced by using $K$.
- An attacker shouldn’t be able to generate a valid $(m, \text{MAC}_K(m))$, even after seeing many valid message-MAC pairs.
- It aims to protect against undetected modifications on messages, not the contents.
  - Sender of a message $m$ computes $\text{MAC}_K(m)$ and appends it to the message.
  - Verification: The receiver also computes $\text{MAC}_K(m)$ & compares to the received value.

MACs from Hash Functions

- prefix: $\text{MAC}_K(m) = H(K || m)$
  - not secure; extension attack.
- suffix: $\text{MAC}_K(m) = H(m || K)$
  - mostly ok; problematic if $H$ is not collision resistant.
- send half of the digest
- envelope: $\text{MAC}_K(m) = H(K_1 || m || K_2)$
- HMAC: $\text{MAC}_K(m) = H(K_2 || H(K_1 || m))$
  - provably secure; popular in Internet standards.