

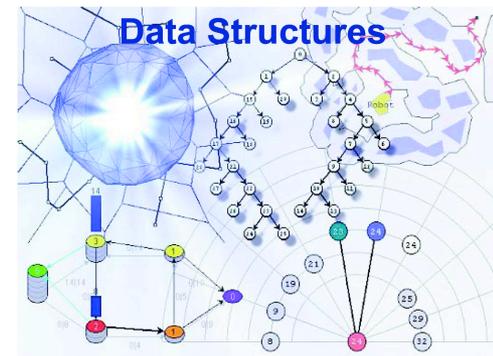
BBM 201

DATA STRUCTURES

Lecture 6: Stacks and Queues



2015-2016 Fall





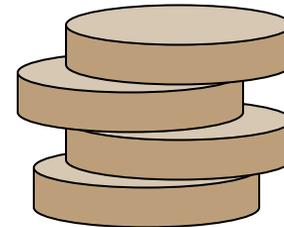
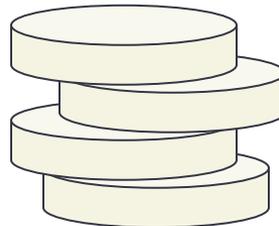
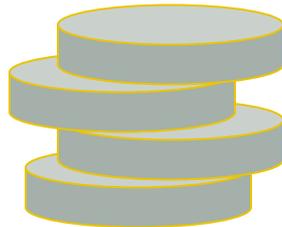
6		9
13	15	60
96	²⁹¹ _{2.94}	297



Stacks

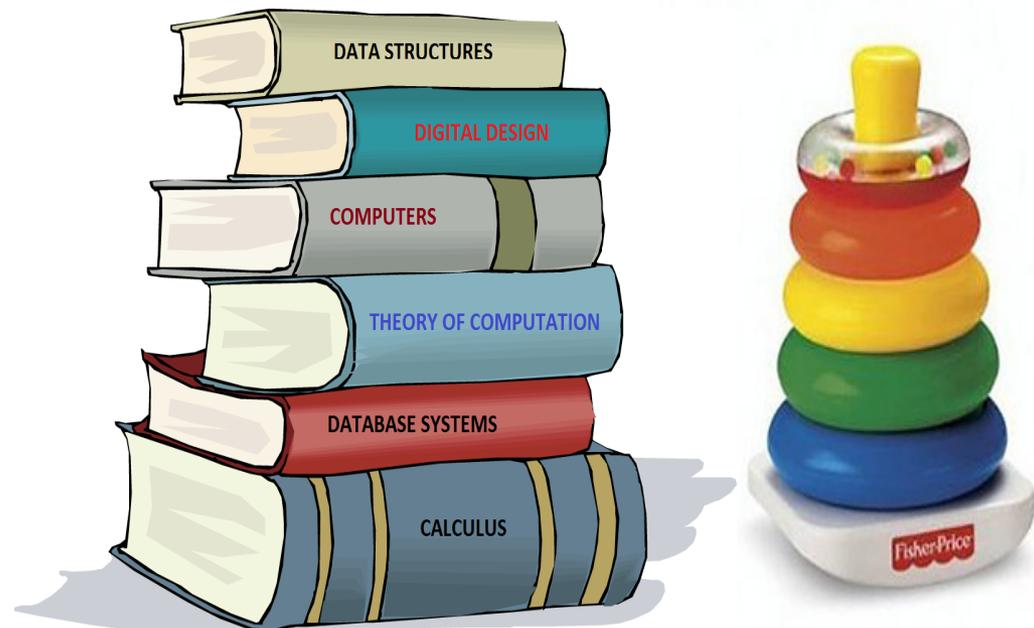
- A list on which insertion and deletion can be performed.
 - Based on Last-in-First-out (LIFO)
- Stacks are used for a number of applications:
 - Converting a decimal number into binary
 - Program execution
 - Parsing
 - Evaluating postfix expressions
 - Towers of Hanoi

...

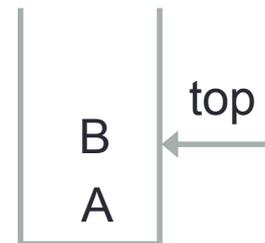
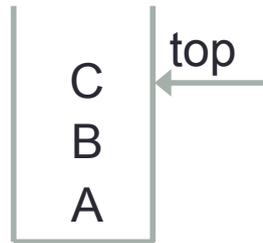
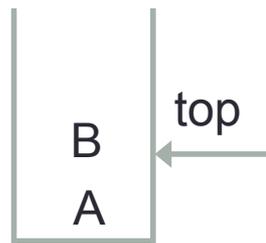
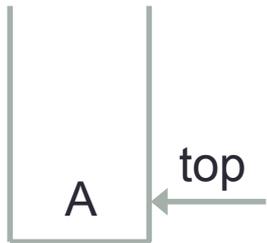


Stacks

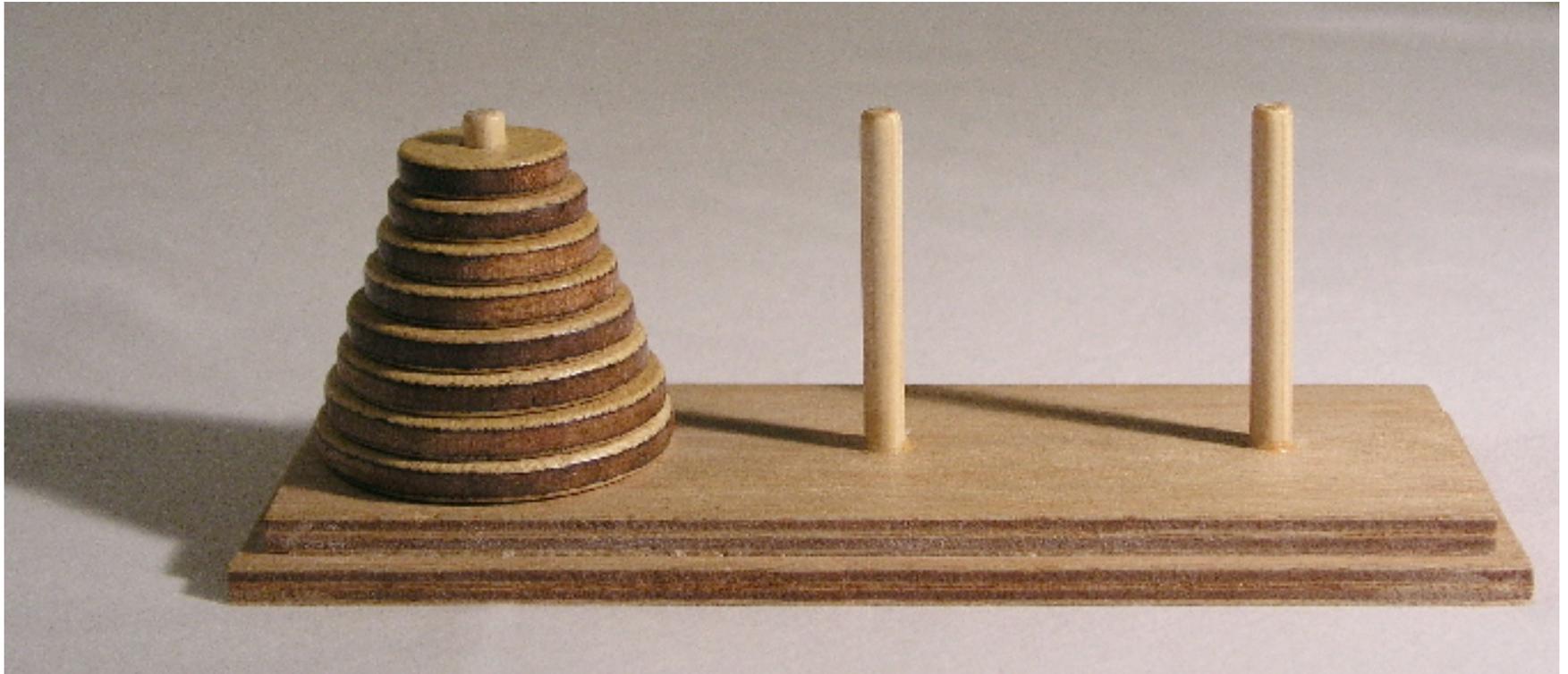
A stack is an ordered lists in which insertions and deletions are made at one end called the **top**.



Stacks

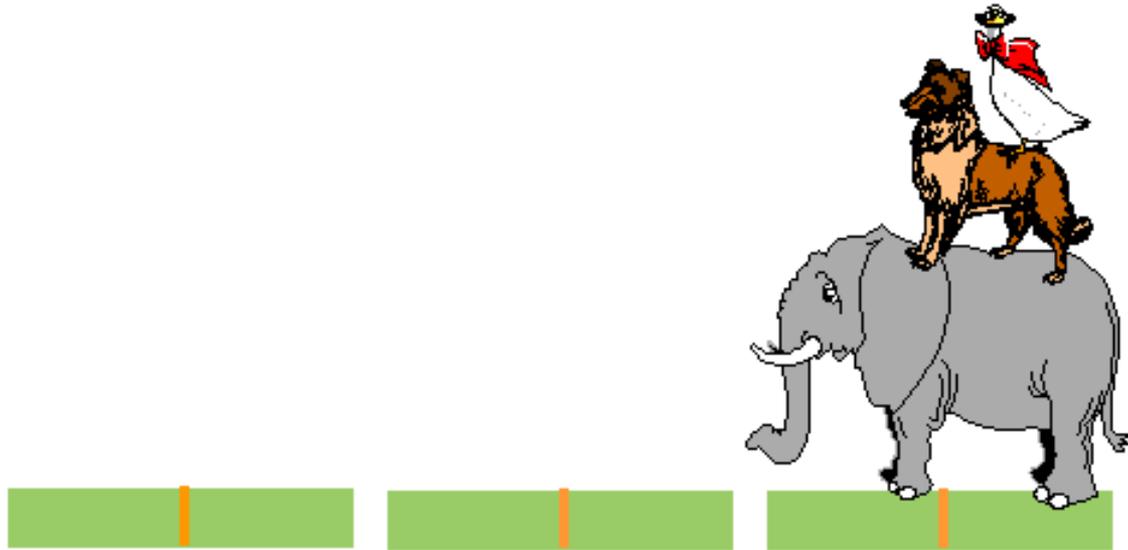


Towers of Hanoi

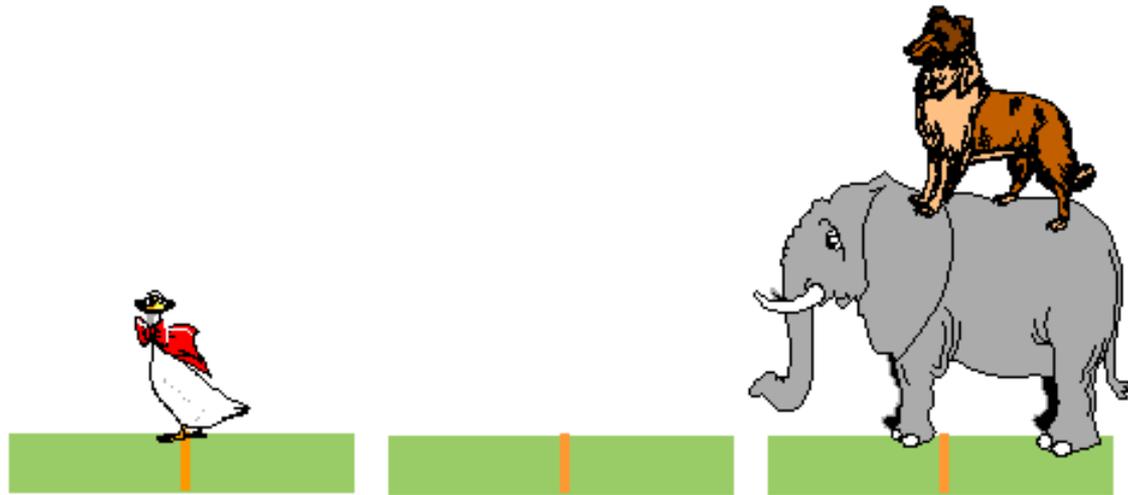


Object of the game is to move all the disks (animals) over to Tower 3. But you cannot place a larger disk onto a smaller disk.

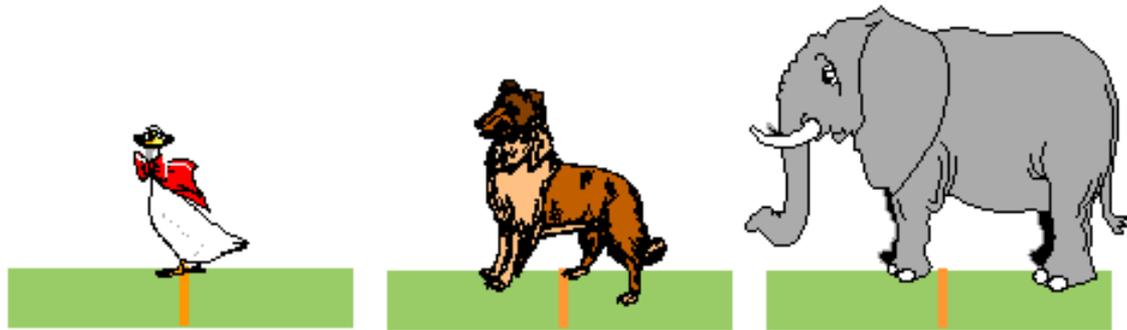
Towers of Hanoi



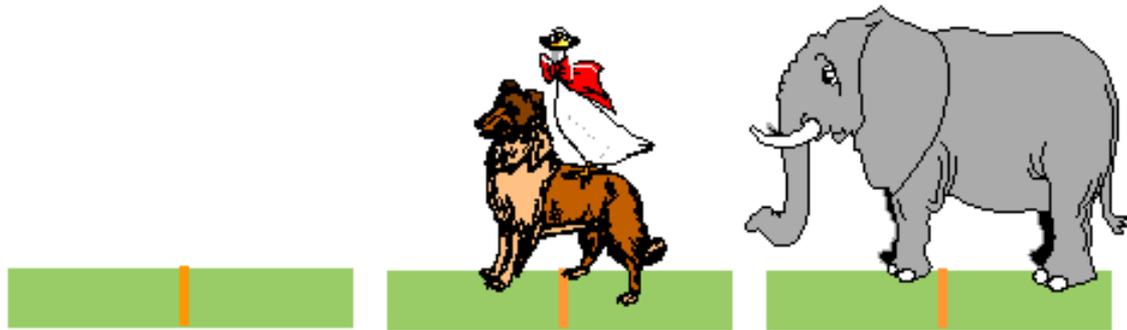
Towers of Hanoi



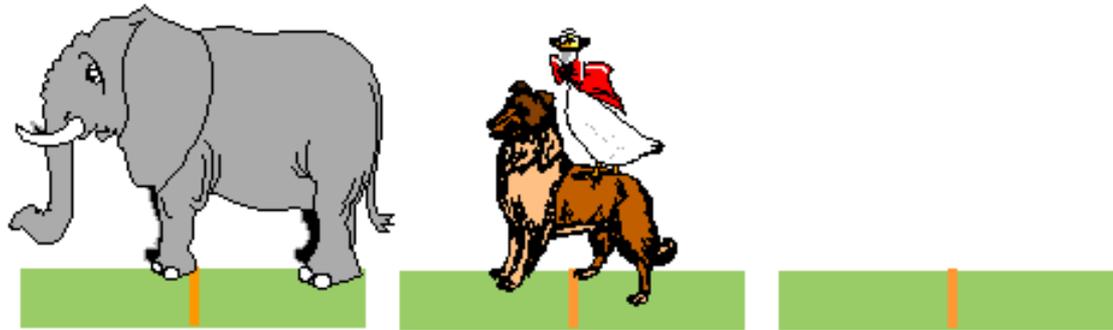
Towers of Hanoi



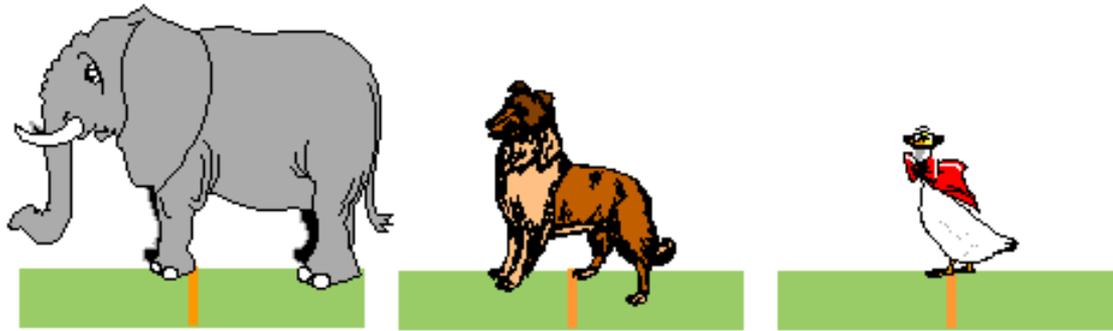
Towers of Hanoi



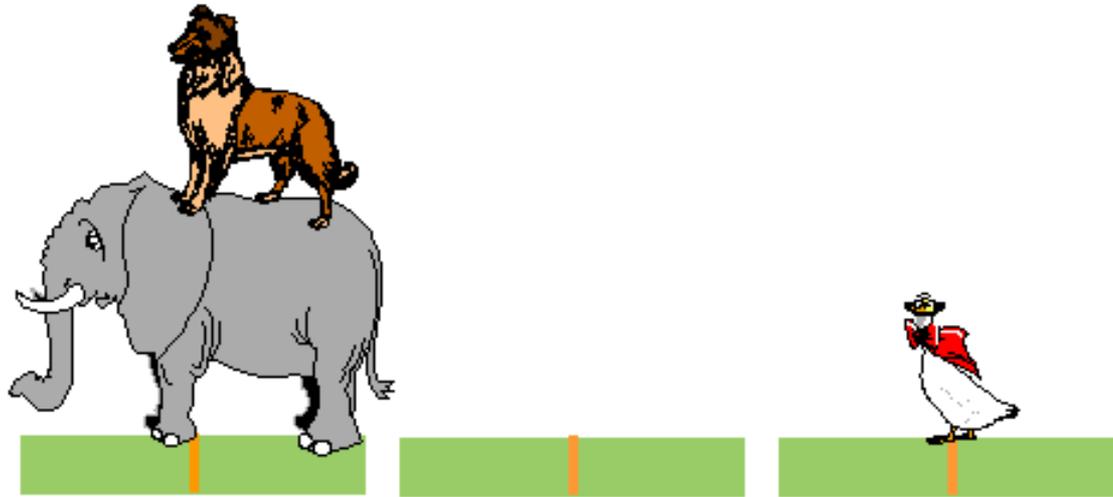
Towers of Hanoi



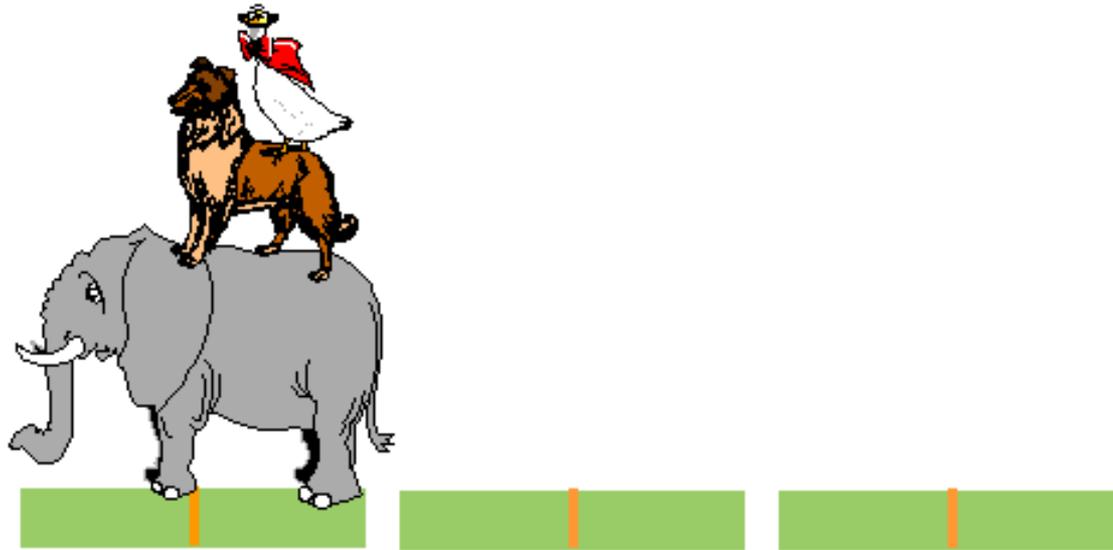
Towers of Hanoi



Towers of Hanoi



Towers of Hanoi

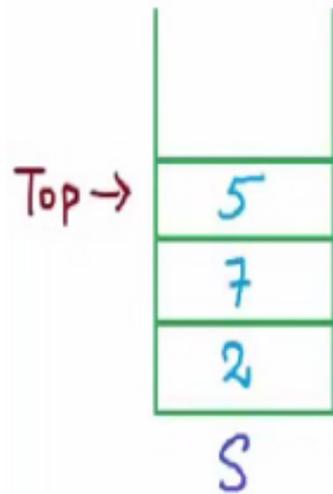


Stack Operations

1. Pop()
 2. Push(x)
 3. Top()
 4. IsEmpty()
- An insertion (of, say x) is called **push** operation and removing the most recent element from stack is called **pop** operation.
 - **Top** returns the element at the top of the stack.
 - **IsEmpty** returns true if the stack is empty, otherwise returns false.

All of these take constant time - $O(1)$

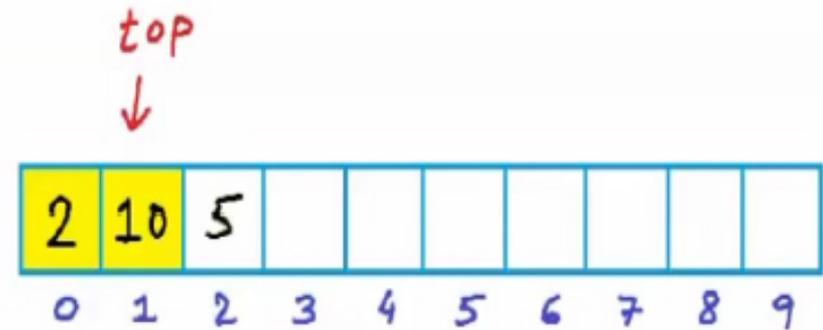
Example



- Push(2)
- Push(10)
- Pop()
- Push(7)
- Push(5)
- Top(): 5
- IsEmpty(): False

Array implementation of stack (pseudocode)

```
int A[10]
top ← -1 //empty stack
Push(x)
{
    top ← top + 1
    A[top] ← x
}
Pop()
{
    top ← top - 1
}
```



Push(2)

Push(10)

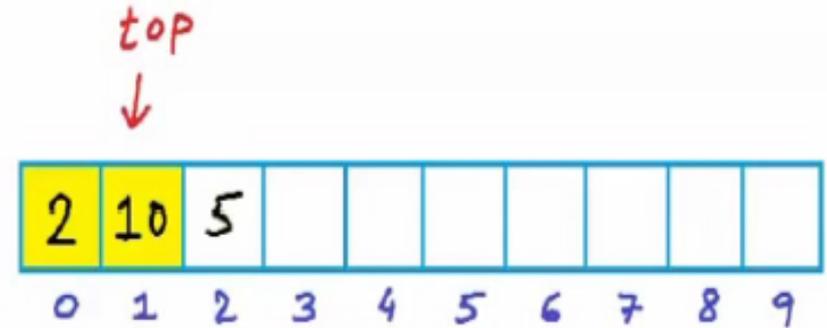
Push(5)

POP()

For an empty stack, top is set to -1.
In push function, we increment top.
In pop, we decrement top by 1.

Array implementation of stack (pseudocode)

```
Top()
{
    return A[top]
}
IsEmpty()
{
    if(top == -1)
        return true
    else
        return false
}
```



Push(2)

Push(10)

Push(5)

POP()

Stack

Data Structure

```
#define MAX_STACK_SIZE 100

typedef struct{
    int deger;
}element;

element stack[MAX_STACK_SIZE];
int top=-1;
```

Push Stack

```
void push(int* top, element item)
{
    if(*top>=MAX_STACK_SIZE){
        isFull();
        return;
    }
    stack[++*top]=item;
}
```

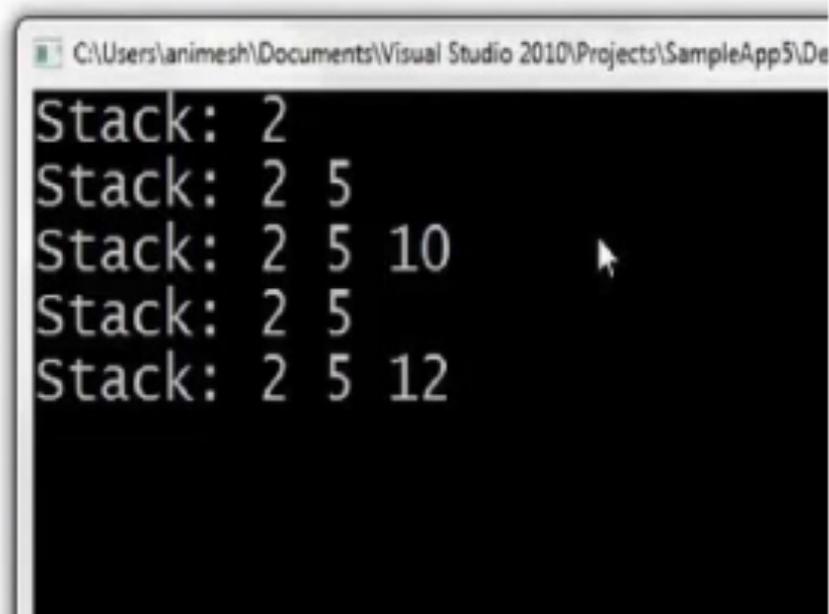
Pop Stack

```
element pop(int* top)
{
    if(*top==-1)
        return empty_stack();
    return stack[(*top)--];
}
```

More array implementation

```
// Stack - Array based implementation.
#include<stdio.h>
#define MAX_SIZE 101
int A[MAX_SIZE];
int top = -1;
void Push(int x) {
    if(top == MAX_SIZE -1) {
        printf("Error: stack overflow\n");
        return;
    }
    A[++top] = x;
}
void Pop() {
    if(top == -1) {
        printf("Error: No element to pop\n");
        return;
    }
    top--;
}
int Top() {
    return A[top];
}
int main() {
}
```

```
void Print() {
    int i;
    printf("Stack: ");
    for(i = 0;i<=top;i++)
        printf("%d ",A[i]);
    printf("\n");
}
int main() {
    Push(2);Print();
    Push(5);Print();
    Push(10);Print();
    Pop();Print();
    Push(12);Print();
}
```



```
C:\Users\animesh\Documents\Visual Studio 2010\Projects\SampleApp5\De
Stack: 2
Stack: 2 5
Stack: 2 5 10
Stack: 2 5
Stack: 2 5 12
```

Check For Balanced Parantheses using Stack

Expression	Balanced?
(A+B)	
{(A+B)+(C+D)}	
{(x+y)*(z)}	
[2*3]+(A)]	
{a+z)	

Check For Balanced Parantheses using Stack

Expression	Balanced?
()	Yes
{()}()	Yes
{()()	No
[]()]	No
{}	No

Need: Count of openings = Count of closings
AND

Any paranthesis opened last should be closed first.

Idea: Create an empty list

- Scan from left to right
 - If opening symbol, add it to the list
 - Push it into the stack
 - If closing symbol, remove last opening symbol of the same type
 - using Pop from the stack
- Should end with an empty list

Check For Balanced Parantheses: Pseudocode

CheckBalancedParanthesis(exp)

{

 n ← length(exp)

 Create a stack: S

 for i ← 0 to n-1

 {

 if exp[i] is '(' or '{' or '['

 Push(exp[i])

 elseif exp[i] is ')' or '}' or ']'

 {if (S is not empty)

 if (top does not pair with exp[i])

 {return false}

 else

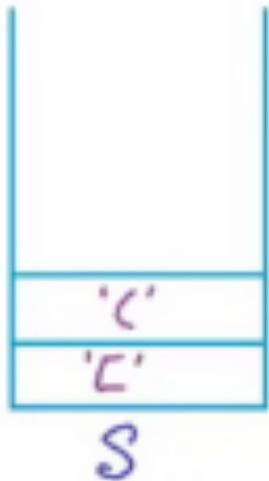
 pop()}}

Return S is empty?

Create a stack of characters and scan this string by using push if the character is an opening parenthesis and by using pop if the character is a closing parenthesis. (See next slide)

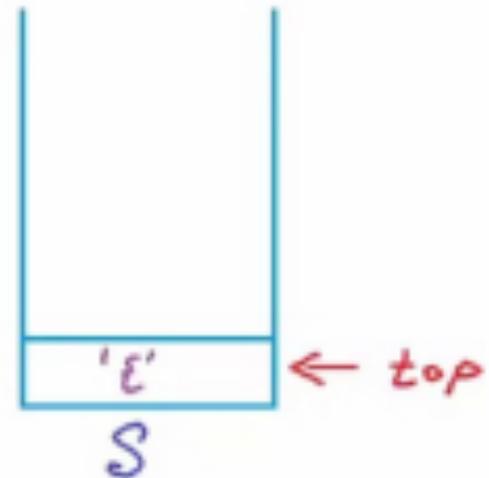
Examples

exp = [(])
 ↑
 i = 2



The pseudo code will return false.

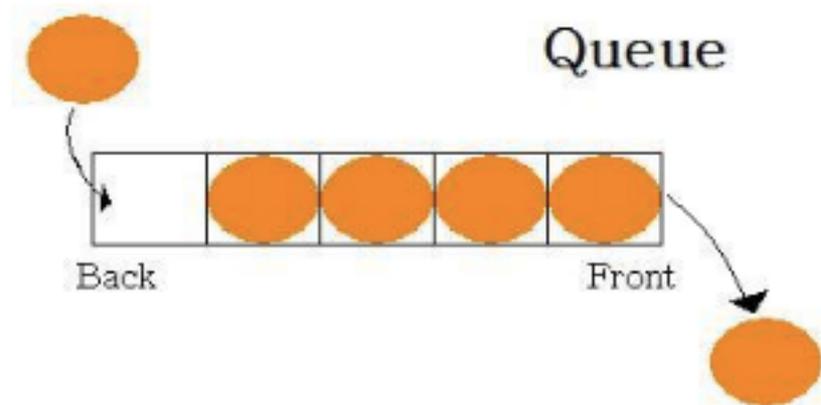
exp = { () () }
 ↑
 i = 5



The pseudo code will return true.

Queues

- A queue is an ordered list on which all insertions take place at one end called the **rear/back** and all deletions take place at the opposite end called the **front**.
 - Based on **First-in-First-out (FIFO)**



Comparison of Queue and Stack

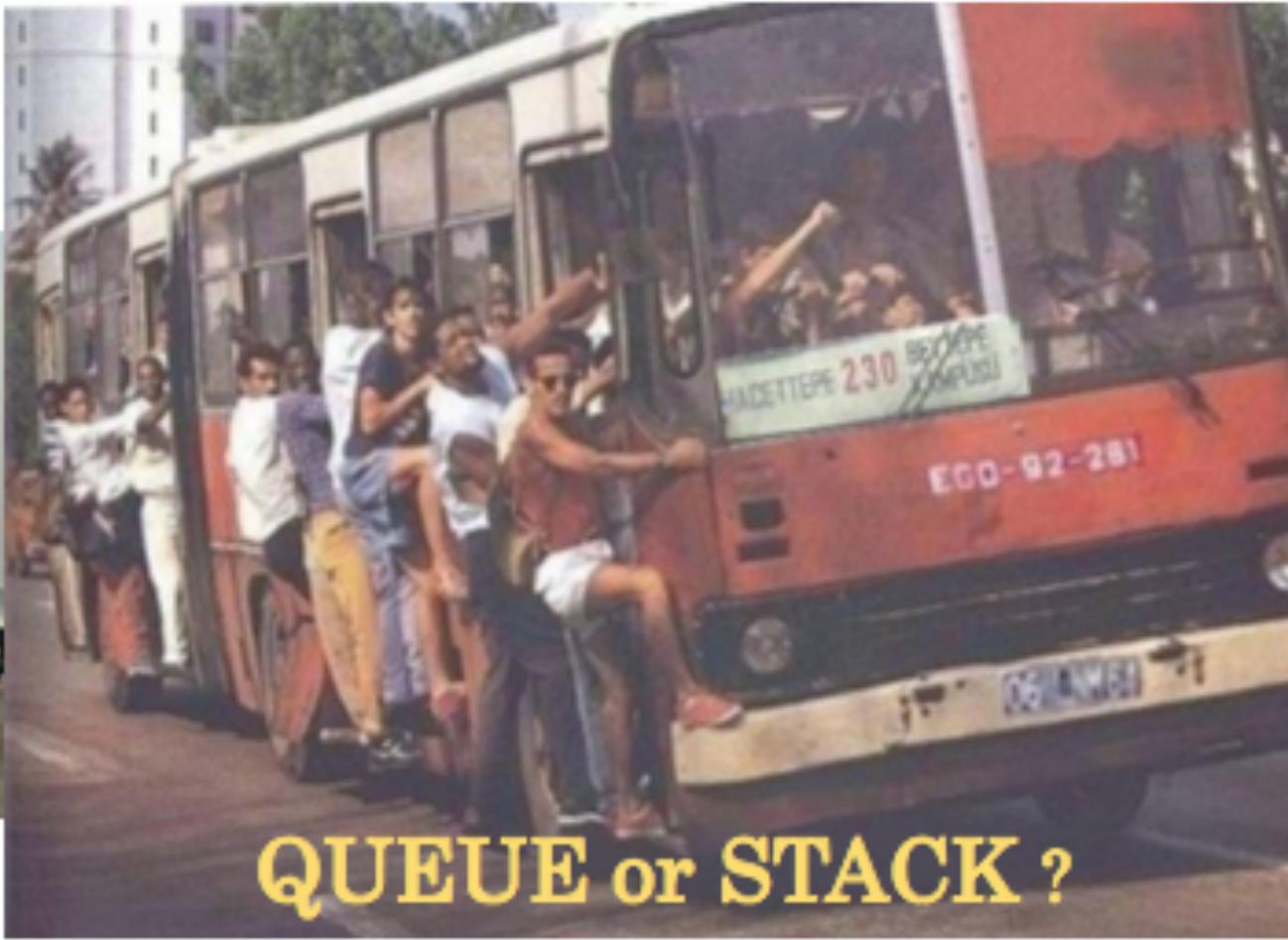
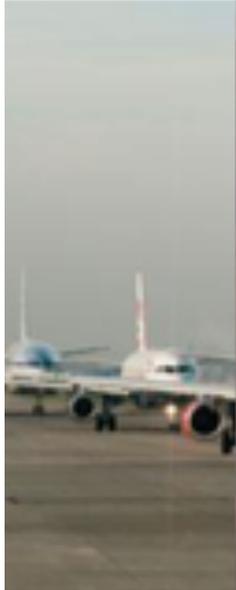
Queue ADT



Queue - First-In-First-Out
(FIFO)



Stack - Last-In-First-Out
(LIFO)



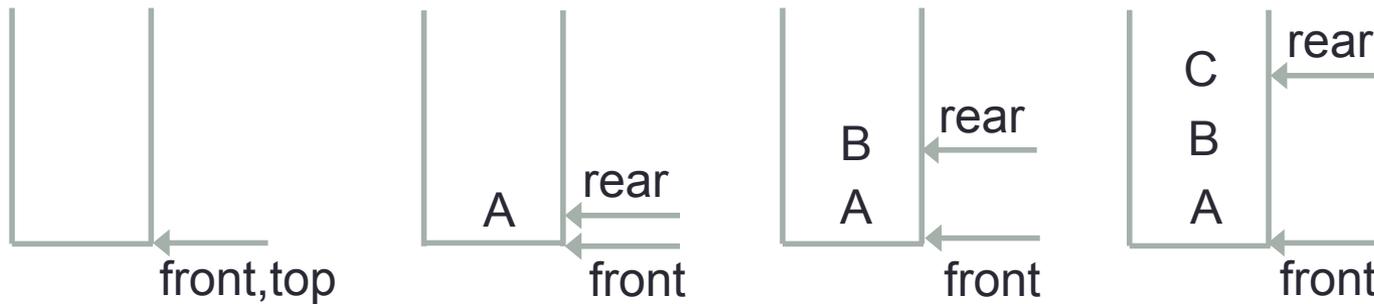
QUEUE or STACK ?



WOWTURK



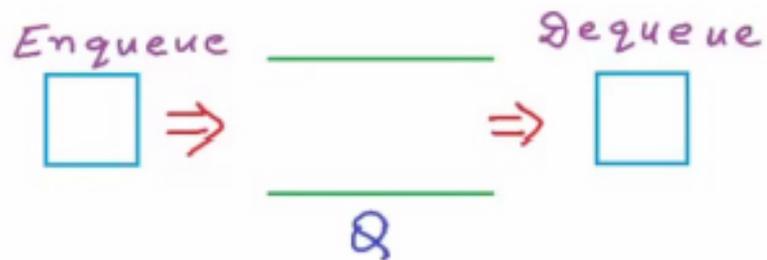
Queues



Queue is a list with the restriction that insertion can be made at one end (**rear**)
And deletion can be made at other end (**front**).

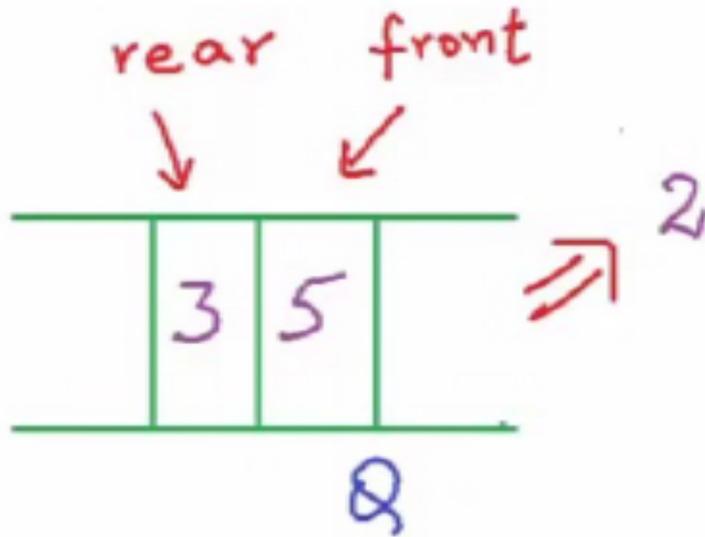
Built-in Operations for Queue

1. Enqueue(x) or Push(x)
2. Dequeue() or Pop()
3. Front(): Returns the element in the front without removing it.
4. IsEmpty(): Returns true or false as an answer.
5. IsFull()



Each operation takes constant time, therefore has $O(1)$ time complexity.

Example



Enqueue(2)

Enqueue(5)

Enqueue(3)

Dequeue() → 2

Front() → 5

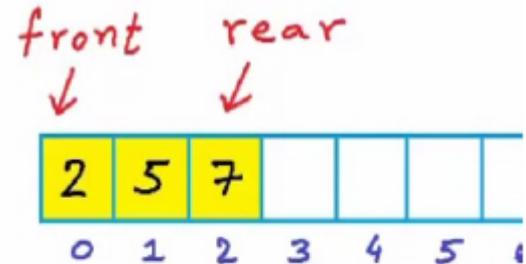
IsEmpty() → False

Applications:

- Printer queue
- Process scheduling

Array implementation of queue (Pseudocode)

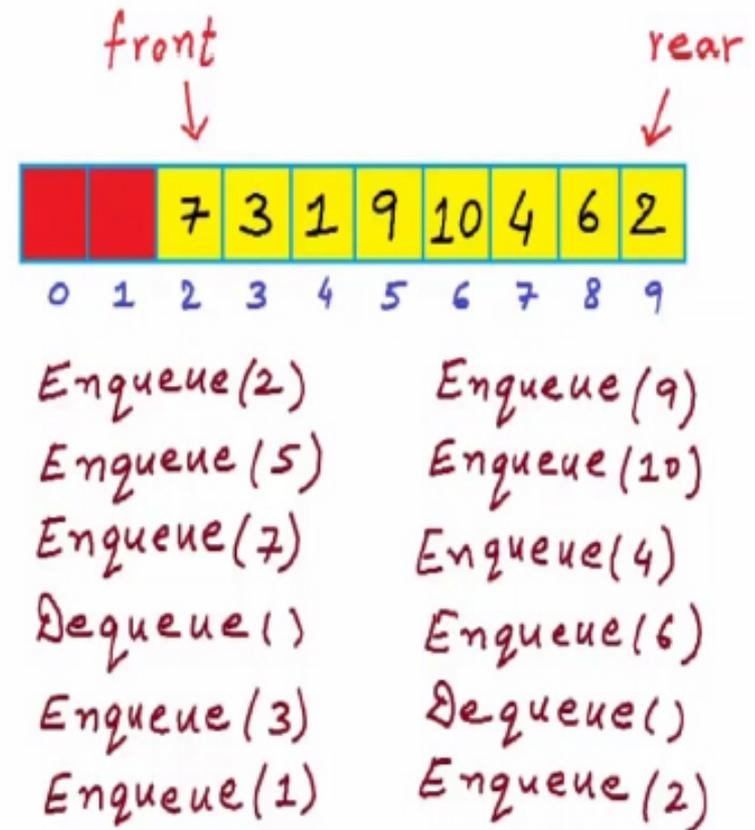
```
int A[10]
front ← -1
rear ← -1
IsEmpty(){
    if (front == -1 && rear == -1)
        return true
    else
        return false}
Enqueue(x){
    if IsFull(){
        return
    }
    elseif IsEmpty(){
        front ← rear ← 0}
    else{
        rear ← rear+1}
    A[rear] ← x}
```



Enqueue(2)
Enqueue(5)
Enqueue(7)

Array implementation of queue (Pseudocode)

```
Dequeue(){  
  if IsEmpty(){  
    return  
  }  
  else if (front == rear){  
    front ← rear ← -1  
  }  
  else{  
    front ← front+1  
  }  
}
```



At this stage, we cannot Enqueue an element anymore.

Queue

Data Structure

```
#define MAX_QUEUE_SIZE 100

typedef struct{
    int deger;
}element;

element queue[MAX_QUEUE_SIZE];
int front=-1;
int rear=-1;
```

Add Queue

```
void addq(int* rear, element item)
{
    if(*rear==MAX_QUEUE_SIZE-1){
        isFull();
        return;
    }
    queue[++*rear]=item;
}
```

Delete Queue

```
element deleteq(int* rear, element item)
{
    if(*front==rear)
        return isEmpty();
    return queue[++*front];
}
```

Circular Queue

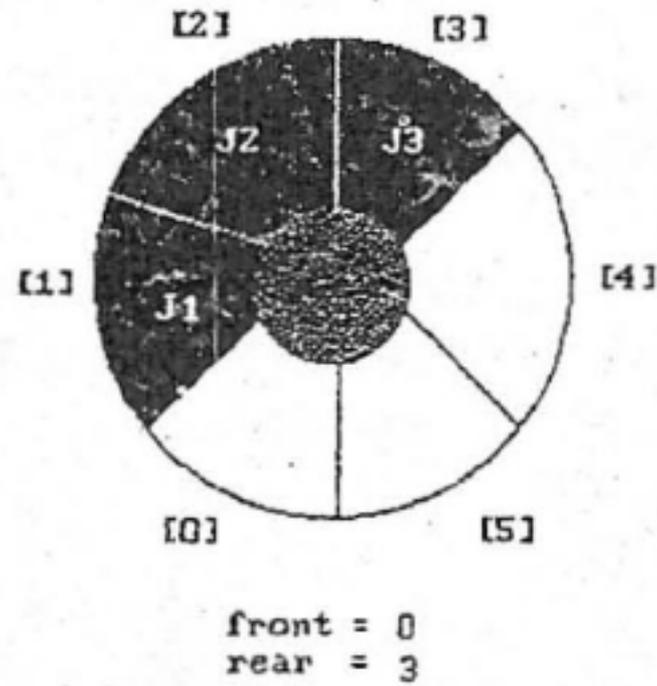
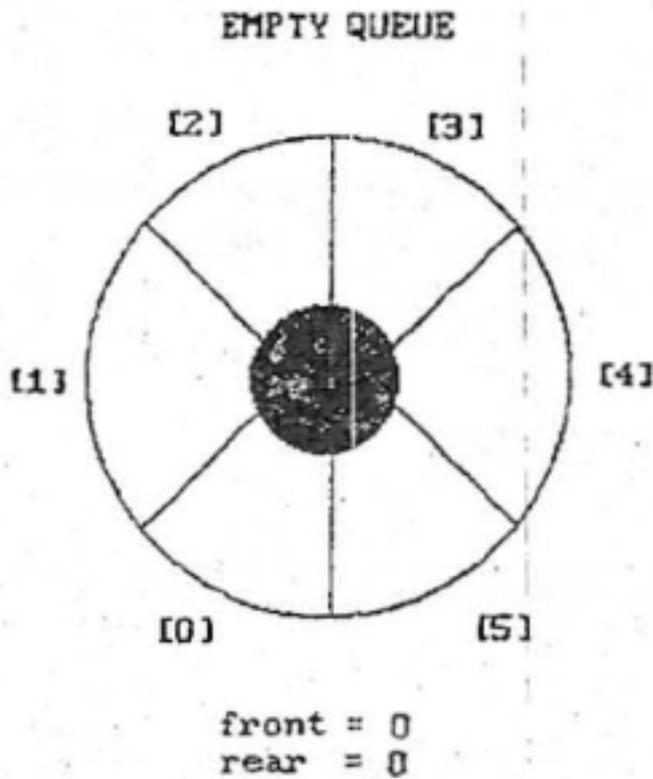
- When the queue is full (the rear index equals to `MAX_QUEUE_SIZE`)
 - We should move the entire queue to the left
 - Recalculate the rear

Shifting an array is time-consuming!

- $O(\text{MAX_QUEUE_SIZE})$

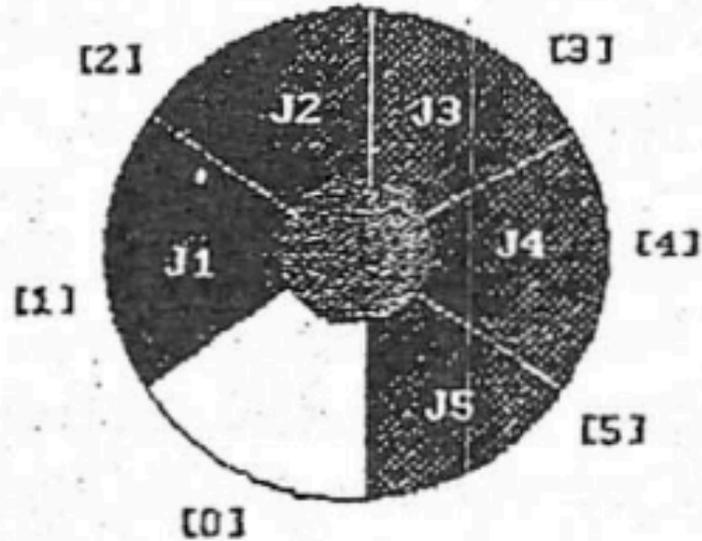
Circular Queue

- More efficient queue representation

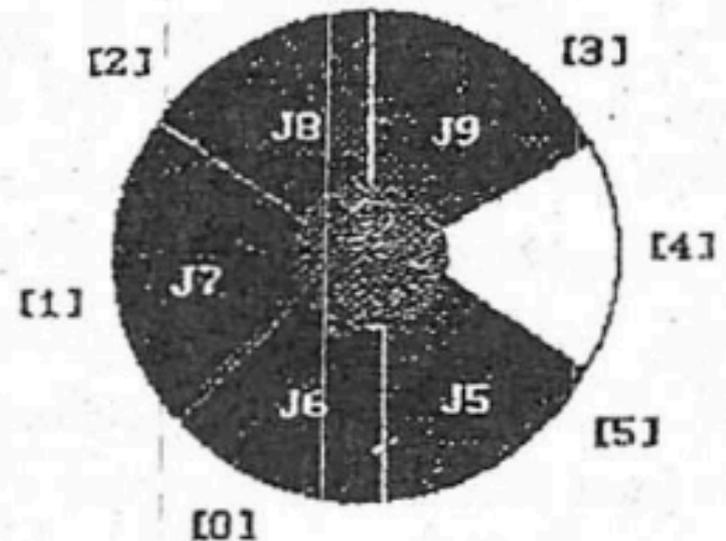


Full Circular Queue

FULL QUEUE



FULL QUEUE



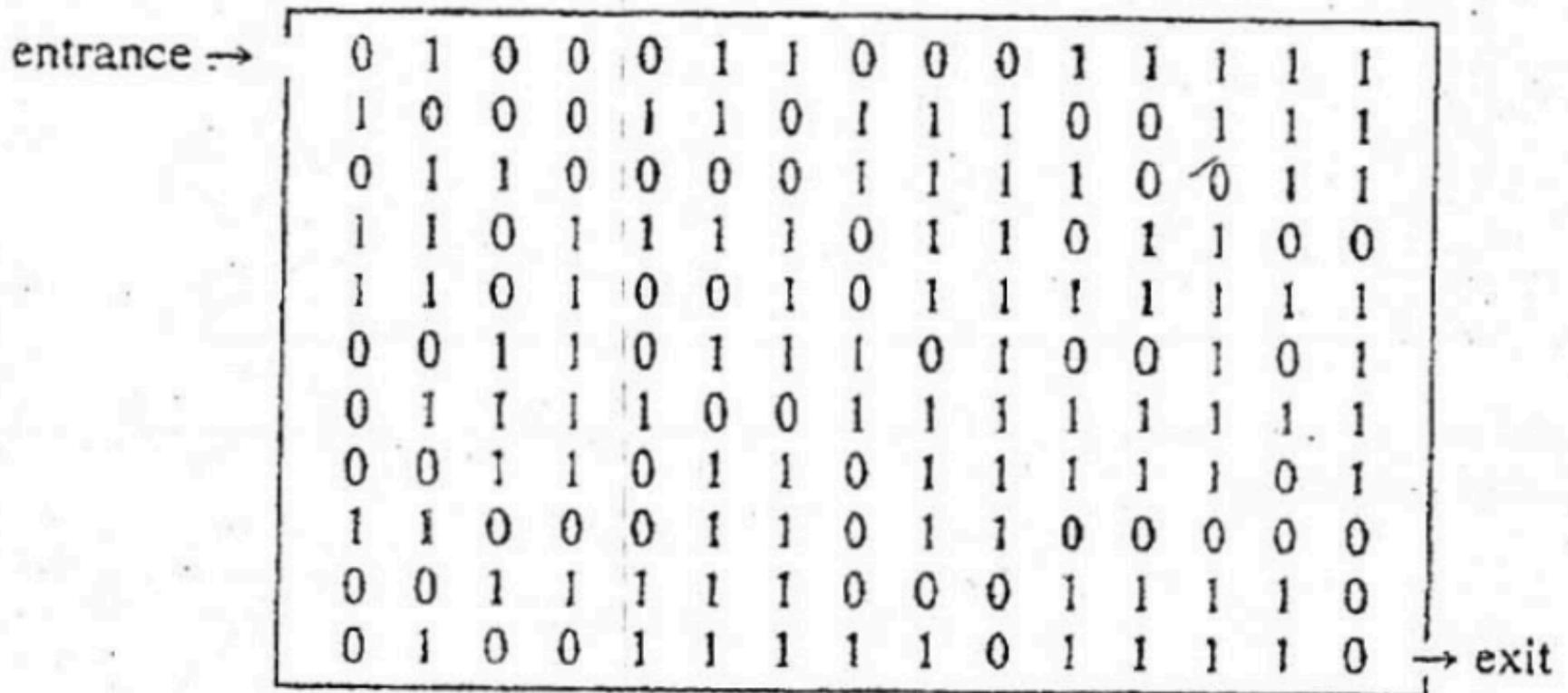
Add Circular Queue

```
void addcircularq(int front, int* rear, element item)
{
    *rear=(*rear+1)%MAX_QUEUE_SIZE;
    if(front==*rear){
        isFull(rear);
        return;
    }
    queue[*rear]=item;
}
```

Delete Circular Queue

```
element deletecircularq(int* front, int arka)
{
    if(*front==rear)
        return isEmpty();
    *front=(*front+1)%MAX_QUEUE_SIZE;
    return queue[*front];
}
```

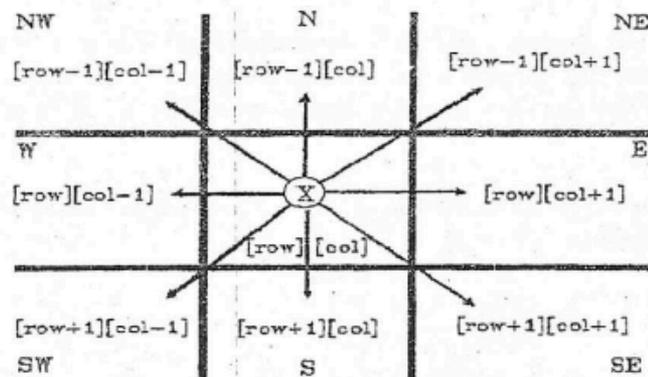
A Mazing Problem



Directions

```
typedef struct{
    short int vert;
    short int horiz;
} offsets;
offsets move[8];
```

Allowable Moves



Name	Dir	$move[dir].vert$	$move[dir].horiz$
N	0	-1	0
NE	1	-1	1
E	2	0	1
SE	3	1	1
S	4	1	0
SW	5	1	-1
W	6	0	-1
NW	7	-1	-1

next_row=row+move[dir].vert;
next_col=col+move[dir].horiz;

IMPLEMENTATION

References

BBM 201 Notes by Mustafa Ege

- <http://www.mycodeschool.com/videos>