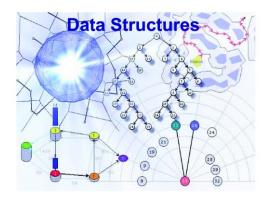
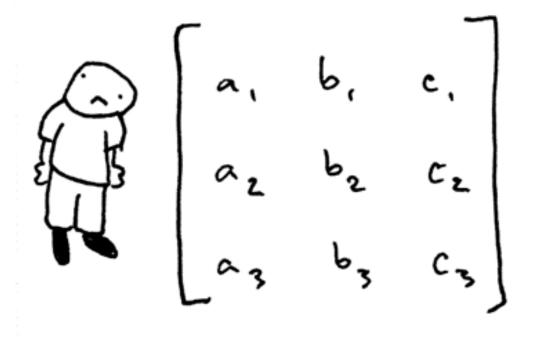
BBM 201 DATA STRUCTURES

Lecture 4:

Lower/Upper Triangular Matrix Band Matrix Sparse Matrix







THE MATRIX!!!!!

Triangular matrix

Upper triangular matrix

Lower triangular matrix

- Does the definition of a special data structure for triangular matrix provide any benefits over a typical matrix in terms of memory and processing time?
- We can insert the items in a single dimensional array:

•	ALT	a ₀₀	a ₁₀	a ₁₁	a ₂₀	a ₂₁	a ₂₂	a ₃₀	a ₃₁	a ₃₂	a ₃₃
		[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]

Number of items in the array becomes:

$$1 + 2 + \dots + (n-1) + n = \frac{n(n+1)}{2}$$

- How can we find the position of u[i][j] in the array?
 - Answer: i=1, there is one item in the 0th row, 2 items in the 1st row.
 - i=2, there is one item in the 0th row, 2 items in the 1st row, 3 items in the 2nd row.
 - Therefore the address of u[i][j] in the array is calculated as below:

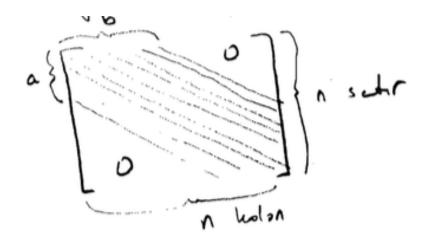
$$k = \mathop{\bigcirc}_{t=0}^{i} (t) + (j) = (0 + 1 + 2 + \dots + i) + (j)$$
$$= \frac{i(i+1)}{2} + (j)$$

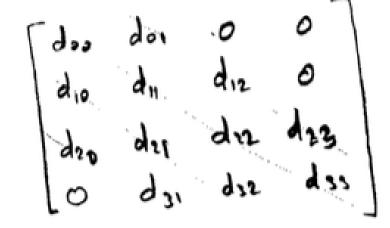
```
void main(void) {
   int alt[MAX_SIZE];
   int i, n;
   scanf("%d", &n); //matrix size
   readtriangularmatrix(alt,n);
   for(i=0; i<=n*(n+1)/2-1; i++)
        printf(" %d", alt[i]);

i=gettriangularmatrix(3,0,n);
   if(i==-2)
        printf("\n invalid index\n");
   else if(i==-1)
        printf("\n access to the upper triangular\n");
   else
        printf("\n the position in 'alt' matrix: %d value: %d \n", i, alt[i]);</pre>
```

$$\begin{bmatrix} B_{11} & B_{12} & 0 & \cdots & \cdots & 0 \\ B_{21} & B_{22} & B_{23} & \ddots & \ddots & \vdots \\ 0 & B_{32} & B_{33} & B_{34} & \ddots & \vdots \\ \vdots & \ddots & B_{43} & B_{44} & B_{45} & 0 \\ \vdots & \ddots & \ddots & B_{54} & B_{55} & B_{56} \\ 0 & \cdots & 0 & B_{65} & B_{66} \end{bmatrix}$$

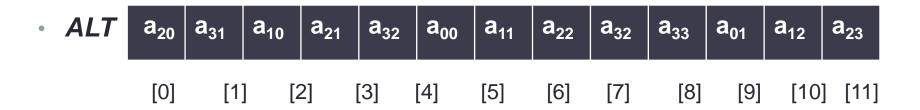
Matrix (n, a): n by n matrix, non-zero entries are confined to a diagonal band, comprising the main diagonal and zero or more diagonals (a-1) on either side.





b=2 a=3 n=4

- What kind of a data structure can we use?
- We can insert the items in a single dimensional array:



What is the number of items in the array?

- What is the number of items in the array?
 - Number of items on and below the diagonal:

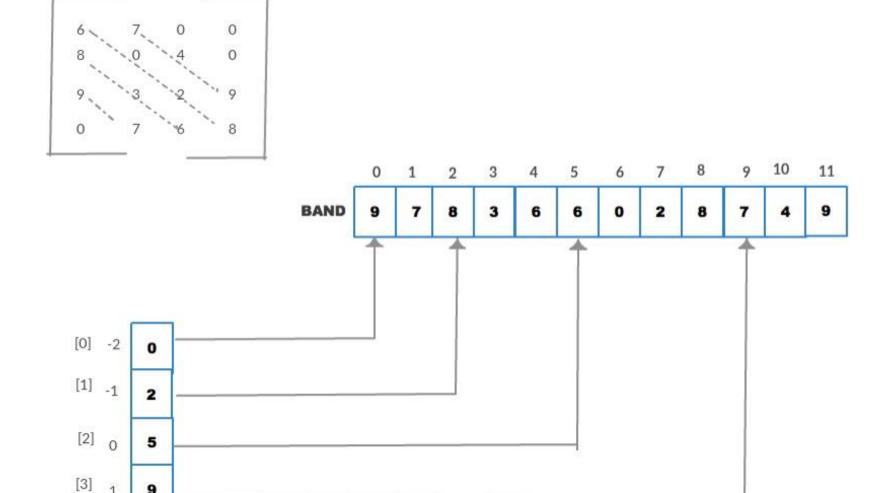
$$n + (n-1) + (n-2) + ... + n - (a-1)$$

Number of items above the diagonal:

$$(n-1)+(n-2)+...+n-(b-1)$$

Sum of these becomes:

$$Sum = n + (n-1) + (n-2) + \dots + n - (a-1) + (n-1) + (n-2) + \dots + n - (b-1)$$
$$= n(a+b-1) - \frac{(a-1)a}{2} - \frac{(b-1)b}{2}$$



```
void main(void) {
   int band[MAX SIZE];
   int search[MAX SIZE];
   int i, n, a, b;
   printf(" n:", &n); scanf("%d", &n);
   printf("a:", &a); scanf("%d", &a);
   printf("b:", &b); scanf("%d", &b);
   buildbandmatrix(band, search, a, b);
   for (i=0; i \le n*(a+b-1)-a*(a-1)/2-b*(b-1)/2-1; i++)
       printf(" %d ",band[i];
    printf("\n");
    for(i=0; i<=a+b-2; i++)
        printf(" %d", search[i]);
    i=getbandmatrix(3,3,n,a,b,search);
    if(i==2)
        printf("\n invalid index");
   else if(i==1)
       printf("\n item to be searched: 0");
    else
        printf("\n item to be searched: %d→%d", i, band[i]);
```

```
void buildbandmatrix(int band[], int search[], int n, int a, int b){
    int i, k, itemnum;
    if(n*(a+b-1)-a*(a-1)/2-b*(b-1)/2 > MAX SIZE) {
        printf("\n not enough memory");
         exit(-1);
     else{
         itemnum=0:
         for(i=-a+1; i<=b-1; i++){ //for each diagonal
              search[i+a-1]=itemnum;
              for(k=0; k<= n-abs(i)-1; k++) //for the current diagonal
                  scanf("%d", &band[search[i+a-1]+k]);
              itemnum = itemnum+(n-abs(i));
```

```
void getbandmatrix(int i, int j, int n, int a, int b, int search[]){
    if(i>=n || i<0 || j>=n || j<0){ //index overflow
          printf("\n invalid index\n");
          return -2;
    else{
          if(j>i) //above the diagonal
             if(j-i<b) //above the upper band
                 return(search[a-1+j-i]+i); //yes
             else //no
                  return -1;
          else if(i-j<a) //below or on the diagonal
                  return(search[j-i+a-1]+j);
          else //not on the band
                  return -1;
```

Sparse Matrix

- Most of the elements are zero.
- It wastes space.

Sparsity: the fraction of zero elements.

Basic matrix operations:

- Creation
- 2. Addition
- 3. Multiplication
- 4. Transpose

A		0	1	2	3	4	5
							_
0		15	0	0	22	0	-15
1		0	11	3	0	0	0
2		0	0	0	-6	0	0
3		0	0	0	0	0	0
4	1	91	0	0	0	0	0
5		0	0	28	0	0	0

Sparse Matrix

Data Structure

```
#define MAX_TERMS 101
typedef struct{
         int col;
         int row;
         int value;
         }term;
term a[MAX_TERMS];
```

- <u>a[0].row:</u> row index
- <u>a[0].col:</u> column index
- <u>a[0].value:</u> number of items in the sparse matrix

Rows and columns are in ascending order!

Sparse Matrix

A		0	1	2	3	4	5	
0		15	0	0	22	0	-15	
1		0	11	3	0	0	0	
2		0	0	0	-6	0	0	
3		0	0	0	0	0	0	
4	1	91	0	0	0	0	0	
5		0	0	28	0	0	0	

	Row	Column	Value
A[0]	6	6	8
A[1]	0	0	15
A[2]	0	3	22
A[3]	0	5	-15
A[4]	1	1	11
A[5]	1	2	3
A[8]	5	2	28

Matrix Transpose

 Replacement of rows and columns in a matrix is called the transpose of the matrix:

$$A = \stackrel{\circ}{e} \stackrel{1}{0} \stackrel{3}{\overset{\circ}{u}} \stackrel{\circ}{u} \qquad \qquad A' = \stackrel{\circ}{e} \stackrel{1}{0} \stackrel{0}{\overset{\circ}{u}} \stackrel{\circ}{u}$$

The item a[i][j] becomes a[j][i].

Matrix Transpose

```
void transpose(term a[],term b[])
  int n,i,j,currentb;
  n=a[0].value; //number of items
  b[0].row=a[0].col; //number of rows
  b[0].col=a[0].row; //number of columns
  b[0].value=n;
  if(n>0){
      currentb=1;
      for(i=0; i<a[0].col; i++)
          for(j=1; j<=n; j++) //find the ones with col i in a
              if(a[j].col==i){
                     b[currentb].row=a[j].col;
                     b[currentb].col=a[j].row;
                     b[currentb].value=a[j].value;
                     currentb++;
```

Question: What is the complexity of this method?

Fast Transpose

```
#define MAX TERM 101
typedef struct{
             int row;
             int col;
             int value;
    } term;
term a[MAX TERM];
void fastTranspose(term a[], term b[])
   int ItemNum[MAX COL], StartPos[MAX COL];
   int i,j,ColNum=a[0].col,TermNum=a[0].value;
   b[0].value=TermNum;
   if(TermNum>0) { //does the item exist?
         for (i=0;i<ColNum;i++)</pre>
               ItemNum[i]=0;
          for (i=1;i<=TermNum;i++)</pre>
               ItemNum[a[i].col]++;
         StartPos[0]=1;
         for (i=1;i<ColNum;i++)</pre>
                StartPos[i]=StartPos[i-1]+ItemNum[i-1];
         for (i=1;i<=TermNum;i++) {</pre>
               j=StartPos[a[i].col]++;
               b[j].row=a[i].col; b[j].col=a[i].row;
               b[j].value=a[i].value;
```

Fast Transpose

- Execute the fastTranspose method.
- Question: What is the complexity of the method?
- Compare its complexity with the previous transpose method.