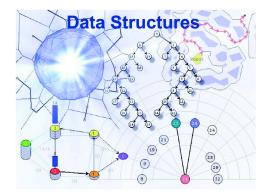
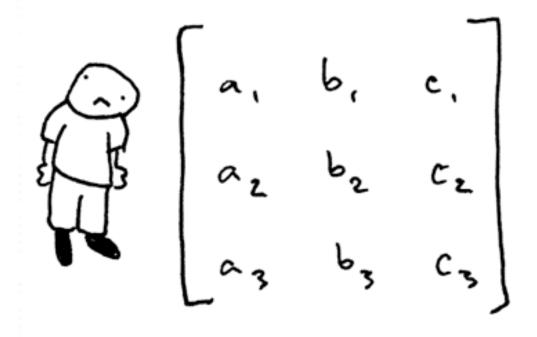
# BBM 201 DATA STRUCTURES

Lecture 4:
Lower/Upper Triangular Matrix
Band Matrix
Sparse Matrix







THE MATRIX!!!!!

#### Triangular matrix

Upper triangular matrix

$$U = \begin{bmatrix} u_{11} & u_{12} & u_{13} & \dots & u_{1n} \\ 0 & u_{22} & u_{23} & & \dots \\ 0 & 0 & u_{33} & & \dots \\ \vdots & & \ddots & & \dots \\ 0 & \ddots & \ddots & \ddots & \dots \\ 0 & \ddots & \ddots & \ddots & \dots \\ 0 & \ddots & \ddots & \ddots & \dots \end{bmatrix}$$

Lower triangular matrix

- Does the definition of a special data structure for triangular matrix provide any benefits over a typical matrix in terms of memory and processing time?
- We can insert the items in a single dimensional array:

•	ALT	a <sub>00</sub>	a <sub>10</sub>	a <sub>11</sub>	a <sub>20</sub>	a <sub>21</sub>	a <sub>22</sub>	a <sub>30</sub>	a <sub>31</sub>	a <sub>32</sub>	a <sub>33</sub>
		[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]

Number of items in the array becomes:

$$1 + 2 + \dots + (n-1) + n = \frac{n(n+1)}{2}$$

$$\begin{bmatrix} a_{00} & 0 & 0 & 0 \\ a_{10} & a_{11} & 0 & 0 \\ a_{20} & a_{21} & a_{22} & 0 \\ a_{30} & a_{31} & a_{32} & a_{33} \end{bmatrix}$$

- How can we find the position of u[i][j] in the array?
  - Answer: i=1, there is one item in the 0<sup>th</sup> row, 2 items in the 1<sup>st</sup> row.
  - i=2, there is one item in the 0<sup>th</sup> row, 2 items in the 1<sup>st</sup> row, 3 items in the 2<sup>nd</sup> row.
  - Therefore the address of u[i][j] in the array is calculated as below:

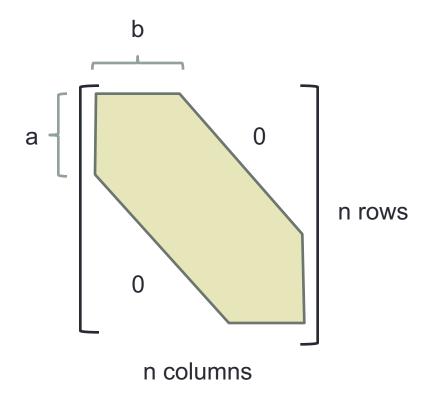
$$k = \sum_{t=1}^{l} (t) + (j) = (1 + 2 + \dots + i) + (j)$$
$$= \frac{i(i+1)}{2} + (j)$$

```
void main(void) {
   int alt[MAX_SIZE];
   int i, n;
   scanf("%d", &n); //matrix size
   readtriangularmatrix(alt,n);
   for(i=0; i<=n*(n+1)/2-1; i++)
        printf("%d", alt[i]);

   i = gettriangularmatrix(3,0,n);
   if(i == -2)
        printf("\n invalid index\n");
   else if(i == -1)
        printf("\n access to the upper triangular\n");
   else
        printf("\n the position in 'alt' matrix: %d value: %d \n", i, alt[i]);
}</pre>
```

$$\begin{bmatrix} B_{11} & B_{12} & 0 & \cdots & \cdots & 0 \\ B_{21} & B_{22} & B_{23} & \ddots & \ddots & \vdots \\ 0 & B_{32} & B_{33} & B_{34} & \ddots & \vdots \\ \vdots & \ddots & B_{43} & B_{44} & B_{45} & 0 \\ \vdots & \ddots & \ddots & B_{54} & B_{55} & B_{56} \\ 0 & \cdots & \cdots & 0 & B_{65} & B_{66} \end{bmatrix}$$

Matrix (n, a): n by n matrix, non-zero entries are confined to a diagonal band, comprising the main diagonal and zero or more diagonals (a-1) on either side.



$$\begin{bmatrix} a_{00} & a_{01} & 0 & 0 \\ a_{10} & a_{11} & a_{12} & 0 \\ a_{20} & a_{21} & a_{22} & a_{23} \\ 0 & a_{31} & a_{32} & a_{33} \end{bmatrix} \qquad b = 2$$

$$a = 3$$

$$n = 4$$

- What kind of a data structure can we use?
- We can insert the items in a single dimensional array:

•	ALT	a <sub>20</sub>	a <sub>31</sub>	a <sub>10</sub>	a <sub>21</sub>	a <sub>32</sub>	a <sub>00</sub>	a <sub>11</sub>	a <sub>22</sub>	a <sub>33</sub>	a <sub>01</sub>	a <sub>12</sub>	a <sub>23</sub>
		[0]	[1]	[2]	[3]	[4]	[5]	[6]	[7]	[8]	[9]	[10]	[11]
		$\begin{bmatrix} a_{00} \\ a_{10} \\ a_{20} \\ 0 \end{bmatrix}$	$a_{01} \\ a_{11} \\ a_{21} \\ a_{31}$	$0 \\ a_{12} \\ a_{22} \\ a_{32}$	$\begin{bmatrix} 0 \\ 0 \\ a_{23} \\ a_{33} \end{bmatrix}$	a b n	= 3 = 2 = 4						

What is the number of items in the array?

- What is the number of items in the array?
  - Number of items on and below the diagonal:

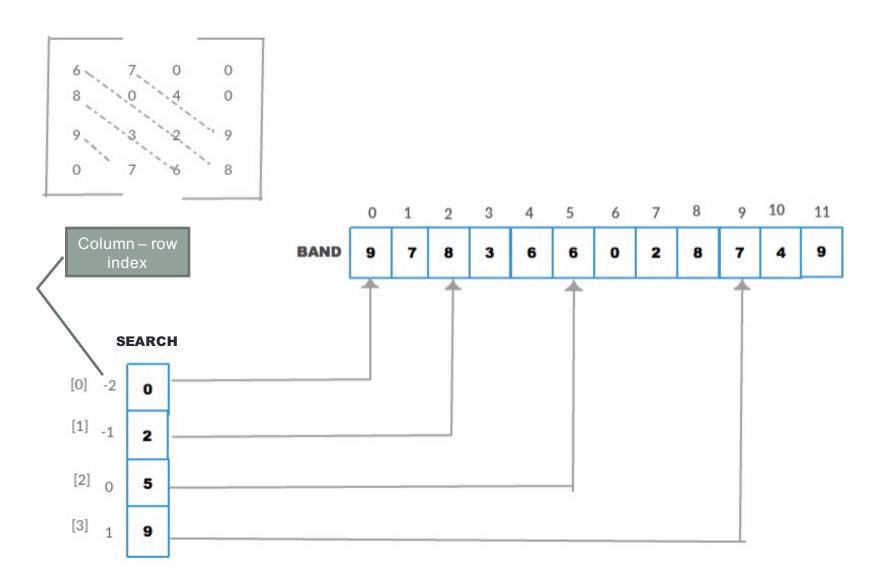
$$n + (n-1) + (n-2) + ... + n - (a-1)$$

Number of items above the diagonal:

$$(n-1)+(n-2)+...+n-(b-1)$$

Sum of these becomes:

$$Sum = n + (n-1) + (n-2) + \dots + n - (a-1) + (n-1) + (n-2) + \dots + n - (b-1)$$
$$= n(a+b-1) - \frac{(a-1)a}{2} - \frac{(b-1)b}{2}$$



```
void main(void){
   int band[MAX SIZE];
   int search[MAX SIZE];
  int i, n, a, b;
  printf(" n:", &n); scanf("%d", &n);
  printf(" a:", &a); scanf("%d", &a);
  printf(" b:", &b); scanf("%d", &b);
  buildbandmatrix(band, search, a, b);
   for (i=0; i \le n*(a+b-1) - a*(a-1)/2 - b*(b-1)/2 - 1; i++)
     printf(" %d ",band[i];
  printf("\n");
  for(i=0; i <= a+b-2; i++)
     printf(" %d", search[i]);
   i = getbandmatrix(3, 3, n, a, b, search);
  if(i == -2)
     printf("\n invalid index");
   else if(i == -1)
     printf("\n item to be searched: 0");
   else
      printf("\n item to be searched: %d -> %d", i, band[i]);
```

```
void buildbandmatrix(int band[], int search[], int n, int a, int b){
   int i, k, itemnum;
   if (n*(a+b-1) - a*(a-1)/2 - b*(b-1)/2 > MAX SIZE){
      printf("\n not enough memory");
      exit(-1);
   else{
      itemnum = 0;
      for (i=-a+1; i \le b-1; i++) \{ //for each diagonal \}
         search[i+a-1] = itemnum;
         for (k=0; k \le n-abs(i)-1; k++) //for the current diagonal
            scanf("%d", &band[search[i+a-1]+k]);
         itemnum = itemnum+(n-abs(i));
```

```
void getbandmatrix(int i, int j, int n, int a, int b, int search[]){
  if(i>=n || i<0 || j>=n || j<0) //index overflow
     printf("\n invalid index\n");
     return -2;
  else
     if(j>i) //above the diagonal
         if(j-i<b) //above the upper band
           return(search[a-1+j-i]+i); //yes
        else //no
           return -1;
     else if (i-j<a) //below or on the diagonal
         return(search[j-i+a-1]+j);
     else //not on the band
        return -1;
```

# **Sparse Matrix**

- Most of the elements are zero.
- It wastes space.

**Sparsity:** the fraction of zero elements.

#### Basic matrix operations:

- 1. Creation
- 2. Addition
- 3. Multiplication
- 4. Transpose

A		0	1	2	3	4	5
0		15	0	0	22	0	-15
1		0	11	3	0	0	0
2		0	0	0	-6	0	0
3		0	0	0	0	0	0
4	1	91	0	0	0	0	0
5		0	0	28	0	0	0

## **Sparse Matrix**

#### **Data Structure**

```
#define MAX_TERMS 101
typedef struct{
          int col;
          int row;
          int value;
        } term;
term a[MAX_TERMS];
```

- <u>a[0].row:</u> row index
- <u>a[0].col:</u> column index
- <u>a[0].value:</u> number of items in the sparse matrix

Rows and columns are in ascending order!

# **Sparse Matrix**

A		0	1	2	3	4	5
0		15	0	0	22	0	-15
1		0	11	3	0	0	0
2		0	0	0	-6	0	0
3		0	0	0	0	0	0
4	1	91	0	0	0	0	0
5		0	0	28	0	0	0

	Row	Column	Value
A[0]	6	6	8
A[1]	0	0	15
A[2]	0	3	22
A[3]	0	5	-15
A[4]	1	1	11
A[5]	1	2	3
A[8]	5	2	28

# **Matrix Transpose**

 Replacement of rows and columns in a matrix is called the transpose of the matrix:

$$A = \begin{bmatrix} 1 & 3 \\ 0 & 4 \end{bmatrix} \qquad A^T = \begin{bmatrix} 1 & 0 \\ 3 & 4 \end{bmatrix}$$

The item a[i][j] becomes a[j][i].

# **Matrix Transpose**

```
void transpose(term a[],term b[])
  int n,i,j,currentb;
  n=a[0].value; //number of items
 b[0].row=a[0].col; //number of rows
 b[0].col=a[0].row; //number of columns
 b[0].value=n;
  if(n>0){
    currentb=1;
    for(i=0; i<a[0].col; i++)
      for(j=1; j<=n; j++) //find the ones with col i in a
        if(a[i].col==i){
          b[currentb].row = a[j].col;
          b[currentb].col = a[j].row;
          b[currentb].value = a[j].value;
          currentb++;
```

Question: What is the complexity of this method?

## **Fast Transpose**

```
#define MAX TERM 101
typedef struct{
            int row;
            int col;
            int value;
              } term;
term a [MAX TERM];
void fastTranspose(term a[], term b[])
  int ItemNum[MAX COL], StartPos[MAX COL];
  int i, j, ColNum=a[0].col, TermNum=a[0].value;
  b[0].value=TermNum;
  if(TermNum>0){ //does the item exist?
    for (i = 0; i < ColNum; i++)
      ItemNum[i] = 0;
    for(i = 1; i <= TermNum; i++)
      ItemNum[a[i].col]++;
    StartPos[0] = 1;
    for (i = 1; i < ColNum; i++)
      StartPos[i] = StartPos[i-1] + ItemNum[i-1];
    for(i = 1; i <= TermNum; i++) {
      j = StartPos[a[i].col]++;
      b[j].row = a[i].col;
      b[i].col = a[i].row;
      b[j].value = a[i].value;
```

# **Fast Transpose**

- Execute the fastTranspose method.
- Question: What is the complexity of the method?
- Compare its complexity with the previous transpose method.