

# Recursions

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## Plan

1. Convergence of sequences
2. Fractals
3. Counting binary trees

## Convergence of Sequences

In the previous lecture we considered a continued fraction for  $\sqrt{2}$ :

$$\sqrt{2} = 1 + \frac{1}{2 + \frac{1}{2 + \frac{1}{2 + \frac{1}{2 + \dots}}}}$$

This fraction can be written in a recursive form

$$x_{n+1} = \frac{1}{x_n + 2}$$

$$x_0 = 0$$

Here are a few first values of the above sequence (coded in *Mathematica*)

```
x[0] = 0; x[n_] :=  $\frac{1}{x[n-1] + 2}$  ;
Table[x[n], {n, 0, 10}] // N
{0., 1., 0.333333, 0.428571, 0.411765, 0.414634,
 0.414141, 0.414226, 0.414211, 0.414214, 0.414213}
```

This numeric experiment suggests that

$$|x_{n+1} - x_n| \rightarrow 0$$

Therefore, we say that a given sequence **converges** if the limit exists:

$$\lim_{n \rightarrow \infty} x_n = a \neq \infty$$

The value to which a sequence converges is called a **fixed point**. For the sequence  $x_n$  a fixed point is  $\sqrt{2} - 1$

$$\lim_{n \rightarrow \infty} x_n = \sqrt{2} - 1$$

We could prove this by the following argument. Let the sequence converges to number  $z$ , namely  $x_n \rightarrow z$ . Clearly, that  $x_{n+1} \rightarrow z$ . We find the value of  $z$  from the sequence definition

$$z = \frac{1}{z + 2}$$

Solving the equation, we get two roots  $z_1 = -\sqrt{2} - 1$ ,  $z_2 = +\sqrt{2} - 1$ . The positive root is the limit.

### ■ Towers of Hanoi

Consider the Towers of Hanoi recursion

$$x_{n+1} = 2x_n + 1$$

$$x_1 = 1$$

Here are the first few values of the sequence

```
Clear[x];
x[1] = 1; x[n_] := 2 x[n - 1] + 1
Table[x[n], {n, 1, 7}] // N
{1., 3., 7., 15., 31., 63., 127.}
```

We say that

$$\lim_{n \rightarrow \infty} x_n = \infty$$

Therefore, the sequence **diverges**.

### ■ More on the Golden Ratio

Consider the following recursion

$$x_{n+1} = \frac{1}{x_n + 1}$$

$$x_0 = 0$$

Does it converge? What does it converge to? Let us assume that  $x_n \rightarrow z$ , then  $x_{n+1} \rightarrow z$  and

$$z = \frac{1}{z+1} \implies z^2 + z - 1 = 0$$

Solving this, yields

$$\lim_{n \rightarrow \infty} x_n = \phi - 1$$

where  $\phi$  is the golden ratio. This immediately leads to a continued fraction for  $\phi$

$$\phi = 1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \frac{1}{1 + \dots}}}}$$

Recall the Euclidean algorithm

$$\begin{aligned} a &= b * q_1 + r_1, & 0 \leq r_1 < b \\ b &= r_1 * q_2 + r_2, & 0 \leq r_2 < r_1 \\ r_1 &= r_2 * q_3 + r_3, & 0 \leq r_3 < r_2 \\ &\dots & \dots \\ r_{k-2} &= r_{k-1} * q_k + r_k, & 0 \leq r_k < r_{k-1} \\ r_{k-1} &= r_k * q_{k+1} + 0 \end{aligned}$$

The continued fraction above implies that all quotients in the Euclidean algorithm applied to  $\text{GCD}(\phi, 1)$  are ones. At the same time, we know that such property holds for  $\text{GCD}(F_{n+1}, F_n)$ . Therefore, we can conject a relation between  $F_n$  and  $\phi$

$$\lim_{n \rightarrow \infty} \frac{F_{n+1}}{F_n} = \phi$$

**Exercise.** Given a sequence

$$x_{n+1} = \sqrt{1 + x_n}$$

$$x_0 = 0$$

that represents a [nested radical](#)

$$\sqrt{1 + \sqrt{1 + \sqrt{1 + \dots}}}$$

What does it converge to?

**Exercise.** Find the fixed point of the following sequence

$$x_{n+1} = \frac{1}{2} \left( x_n + \frac{4}{x_n} \right)$$

$$x_0 = 1$$

Consider a general form recursion

$$x_{n+1} = f(x_n)$$

If  $x_n$  converges to a number  $z^*$ , then  $z^*$  is a fixed point

$$z^* = f(z^*)$$

Solving this equation, we find  $z^*$ . **The functional analysis question:** for what classes of function  $f$  the sequence  $x_n$  converges?

## Fractals

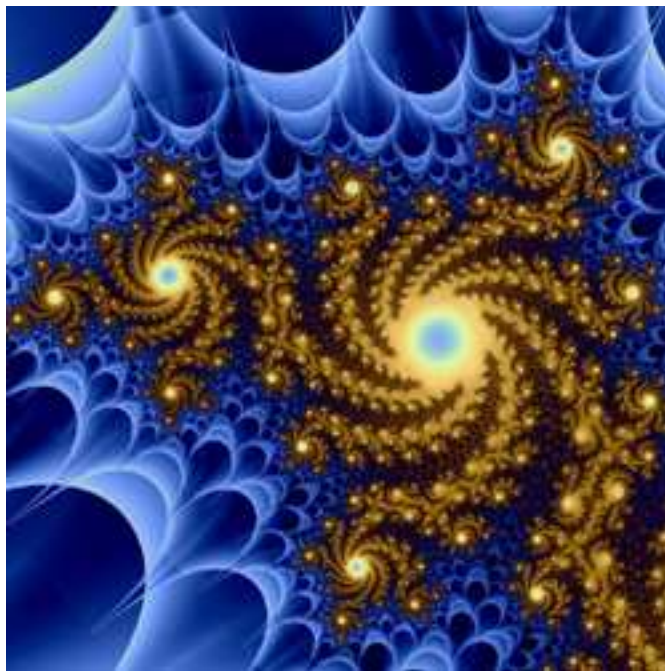
Fractals are geometric objects that are self-similar, i.e. composed of infinitely many pieces, all of which look the same.



Some fractals are mundane



But some fractals are extremely complicated



Since producing fractals requires repeating a certain step over and over again on smaller and smaller scales, it can be easily drawn by a computer.

#### ■ Mandelbrot Set

The [Mandelbrot set](#) is a set of complex numbers  $z$  for which the following recurrence converges

$$a_{n+1} = a_n^2 + z$$

$$a_0 = z$$

Let  $z = 1$ , we get

$$a_{n+1} = a_n^2 + 1$$

$$a_0 = 1$$

Here are the first few value of the sequence

```
Clear[x];
x[0] = 1; x[n_] := x[n - 1]^2 + 1
Table[x[n], {n, 1, 5}] // N
{2., 5., 26., 677., 458330.}
```

The sequence does not converge. However, if  $z = \frac{1}{5}$  then

$$a_0 = \frac{1}{5} = 0.2$$

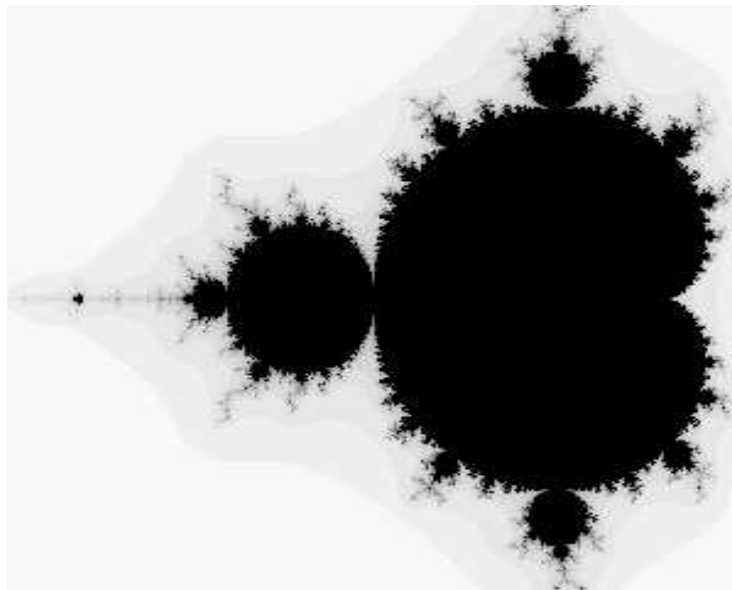
$$a_1 = \left(\frac{1}{5}\right)^2 + \frac{1}{5} = \frac{6}{25} = 0.24$$

$$a_2 = \left(\frac{6}{25}\right)^2 + \frac{1}{5} = \frac{161}{625} = 0.2576$$

$$a_3 = \left(\frac{161}{625}\right)^2 + \frac{1}{5} = \frac{104046}{390625} = 0.26636$$

the sequence does converge to 0.276393.

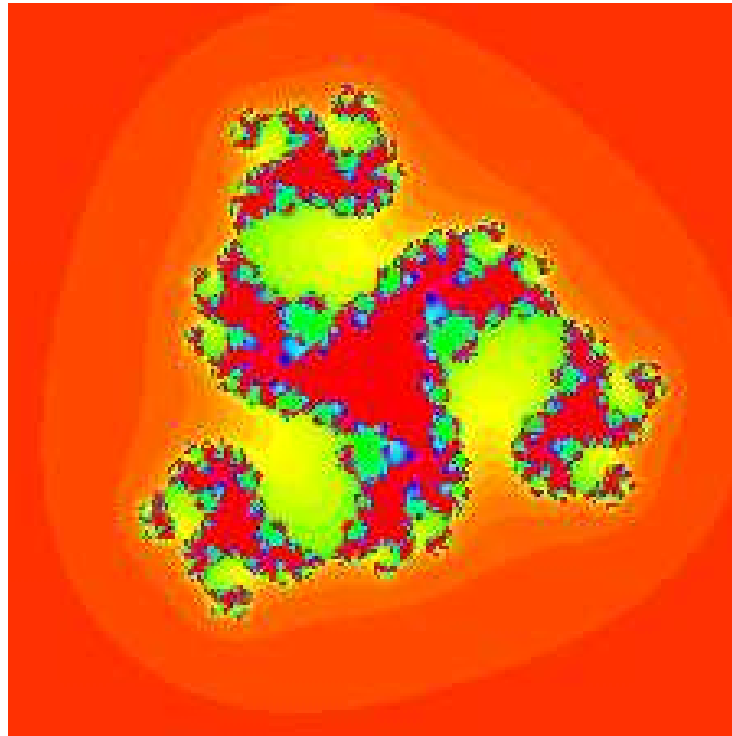
Here is a picture of the number of iterations that takes until a fixed point has been reached. Different shadows of gray correspondent to a different number of iterations.



## ■ Julia Set

**Julia sets** are defined by iterating a function starting at the arbitrary point in the plane. If after some number of iterations the result does not drift to infinity, but instead tends to a fixed point, then that starting point belongs to the Julia set. Here is a picture of the Julia set for

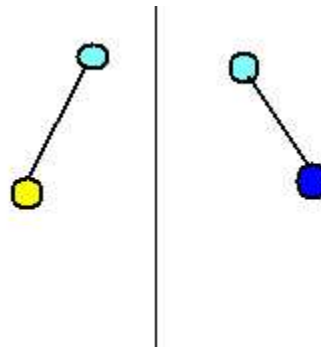
$$x_{n+1} = \text{Conjugate}(x_n)^3 - 0.53 - 0.4 * i$$



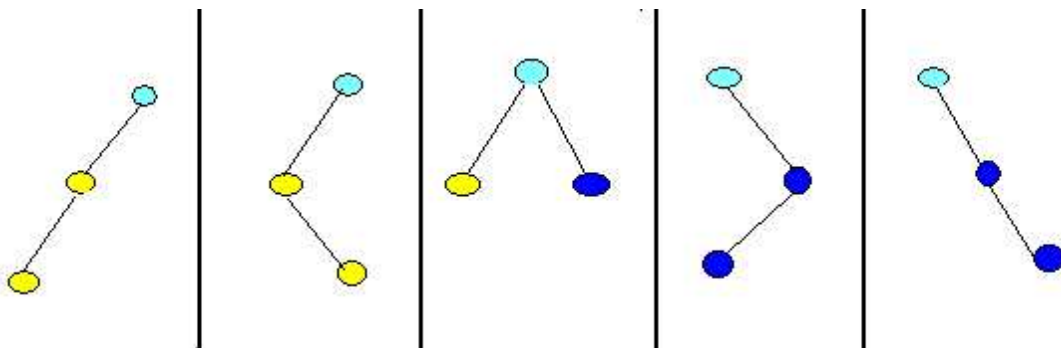
## Counting Binary Trees

A binary tree is made of nodes, where each node contains a "left" reference, a "right" reference, and a data element. The left and right references recursively point to smaller "subtrees" on either side. The topmost node in the tree is called the "root". A recursive definition: a binary tree is either empty or consists of a root, a left subtree and a right subtree.

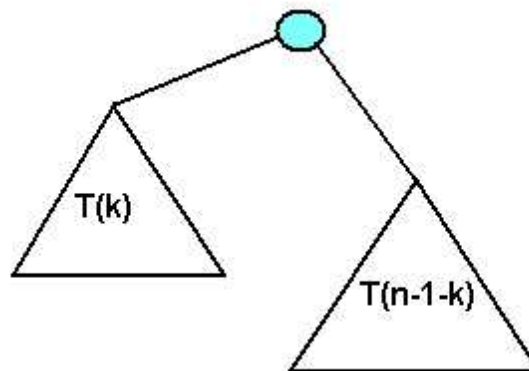
In this section we will count the number of binary trees with  $n$  nodes. Consider a few particular cases. If  $n = 1$ , there is only one binary tree. If  $n = 2$ , there two trees



If  $n = 3$ , there five trees



In general case we will derive a recursive formula for the number of trees based on the recursive definition of a binary tree. Let  $T(n)$  denote the number of binary trees with  $n$  nodes. Suppose the left subtree (LT) has  $k$  nodes, then the right one (RT) has  $n - 1 - k$  nodes.



Thus altogether we can create  $T(k) * T(n - 1 - k)$  binary trees with  $k$  nodes in the left subtree. Since the left subtree can have any number of nodes in the interval  $0 \leq k \leq n - 1$ , we have to sum up over all such  $k$



$$T(n) = \sum_{k=0}^{n-1} T(k) * T(n-1-k)$$

The solution to this recurrence is known as the [Catalan numbers](#) after the Belgian Eugene Charles Catalan:

$$T(n) = \frac{1}{n+1} \binom{2n}{n}$$

where  $\binom{2n}{n}$  stands for binomial coefficients.

**Exercise.** Give a recurrence relation for the number of ways to climb  $n$  stairs if the climber can take either one or two stairs at a time.