

BBM 201 - Data Structures - Fall 2015

Midterm 1

Date: October 26, 2015

Time: 9:30-11:20

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Şube /Section:

Question	1	2	3	4	5	6	7	8	9	10	11	Total
Points	10	8	10	6	10	6	10	8	15	9	8	100
Grade												

1. (10 points) Decide if the following statements are true or false, circle your answer.

The cost of insertion (Push) and deletion (Pop) of an element in a stack is $O(1)$. True False

The ~~complexity~~ ^{order of growth} of the worst-case running time of the 2-Sum problem for an input of size n is $n \log n$. True False

The ~~complexity~~ ^{order of growth} of the worst-case running time of the iterative Fibonacci algorithm for finding f_n is n . True False

The ~~complexity~~ ^{order of growth} of the worst-case running time of binary search for an array of size n is n^2 . True False

$3n^2 + 10n \log n = O(n \log n)$ True False

2. (8 points)

(a) (4 points) Find what RecursiveFunc (given below) does in general.

Finds the maximum of elements in a

(b) (4 points) Print the output of the following program.

```
int RecursiveFunc (int array[], int index, int n)
{
    int val1, val2;
    if ( n==1 )
        return array[index];
    val1 = RecursiveFunc (array, index, n/2);
    val2 = RecursiveFunc (array, index+(n/2), n-(n/2));
    if (val1 > val2)
        return val1;
    else
        return val2;
}
...
int a[8] = 1,2,10,15,16,4,8,2;
printf( "value is % d \n", RecursiveFunc(a,0,8));
```

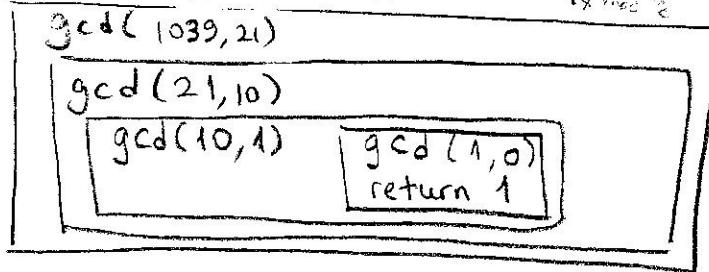
value is 16

3. (10 points) Below is a pseudocode of the Euclidean algorithm that calculates the greatest common divisor (gcd) of any given two integers. Write all recursive calls made to calculate gcd(21, 1039) by using a box diagram.

```

algorithm gcd(x,y)
  if y = 0
    then return(x)
  else return(gcd(y,x mod y))
  
```

$$\begin{array}{l}
 x = 1039 \\
 y = 21
 \end{array}
 \left. \begin{array}{l}
 \begin{array}{r}
 1039 \overline{) 21} = 49 \\
 \underline{-84} \\
 199 - 189 = 10
 \end{array} \\
 x \text{ mod } y
 \end{array} \right\}
 \begin{array}{r}
 21 \overline{) 10} \\
 \underline{20} \\
 1 \text{ x mod } \frac{1}{2}
 \end{array}
 \begin{array}{r}
 10 \overline{) 1} \\
 \underline{10} \\
 0 \text{ x mod } \frac{1}{10}
 \end{array}$$



4. (6 points) Write the complexity of the running time of the following code fragment using big-Oh notation. Show your work.

```

int count = N, i = 0, j = 0;
for (i = 0; i < count; i++)
  if (i % 2 == 0) // i, even
    j++;
  else
    j--;
for (i = count; i > 0; i = i/2)
  for (j = 0; j < i; j++)
    count++;
  
```

$$\begin{aligned}
 & \left\lceil \frac{N}{2} \right\rceil + \left\lceil \frac{N}{2} \right\rceil + N + \frac{N}{2} + \frac{N}{4} + \dots + \frac{N}{2^k} = \\
 & k = \log_2 N \quad = O(N)
 \end{aligned}$$

5. (10 points) How many array accesses do the following code fragments make as a function of N?

(a) (5 points)

```
int count = 0;
for (int i = 0; i < N; i++)
    for (int j = i+1; j < N; j++)
        if (a[i] + a[j] == 0)
            count++;
```

$$2 \cdot [(N-1) + (N-2) + \dots + 1] = 2 \binom{N}{2} = \frac{N(N-1)}{2} \cdot 2$$

(b) (5 points)

```
int count = 0;
for (int i = 0; i < N; i++)
    for (int j = i+1; j < N; j = 2*j)
        if (a[i] + a[j] == 0)
            count++;
```

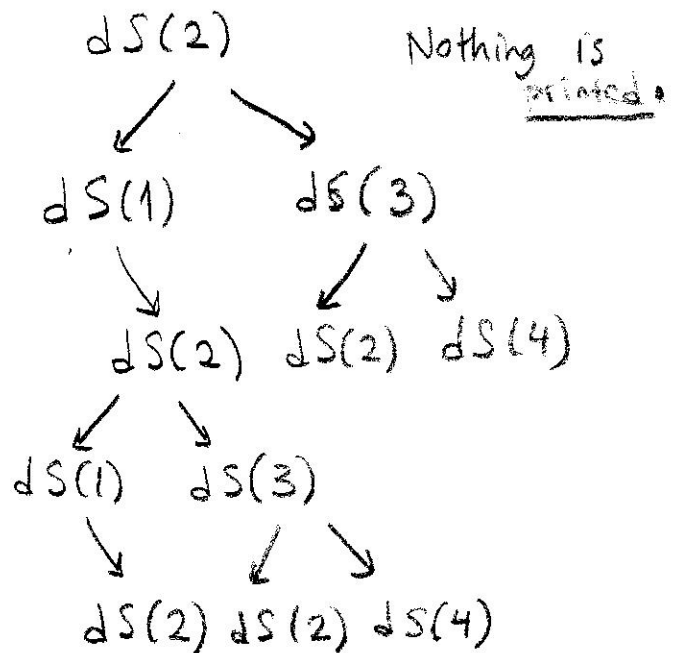
N times

$$2 \left[\log_2(N-1) + \log_2(N-1) + \dots + \log_2(N-1) \right]$$

$\sim 2N \log_2 N$

6. (6 points) What will be the output when doSomething(2) is called?

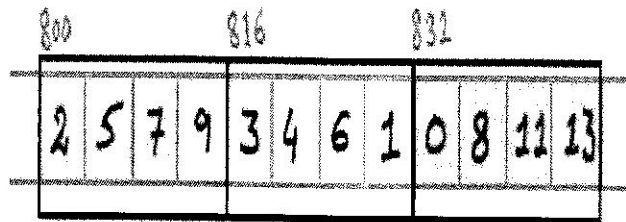
```
void doSomething(int value)
{
    if(0 < value && value < 10)
    {
        doSomething(value - 1);
        doSomething(value + 1);
        printf(" %d", value);
    }
}
```



Never Stops

7. (10 points) If C is the array shown with its address above each node, write what the following lines of a program will print in the empty column.

```
int C[3][2][2];
```



The code	will print:
<code>printf("C=%d", C);</code>	800
<code>printf("C+1=%d", C+1);</code>	816
<code>printf("*(C+1) =%d", *(C+1));</code>	808
<code>printf(" *(C[1]+1)=%d '' , *(C[1]+1)) ;</code>	824
<code>printf(" *(C[1][1]+1)=%d '' , *(C[1][1]+1));</code>	<u>1</u>

8. (8 points) Let an array be defined with `int a[6][8][5][7]`. What is the memory address of `a[3][4][2][1]` if the memory address of `a[0][0][0][0]` is x ? (Assume that size of an integer is 4 bytes.)

$$x + [3 \cdot (8 \cdot 5 \cdot 7) + 4 \cdot (5 \cdot 7) + 2 \cdot 7 + 1] \cdot 4$$

9. (15 points) Let $A[n][n]$ be an upper triangular matrix in the example given below. The elements of this triangular matrix A are stored in the one-dimensional array U as shown.

$$A = \begin{bmatrix} a_{00} & a_{01} & a_{02} & a_{03} \\ 0 & a_{11} & a_{12} & a_{13} \\ 0 & 0 & a_{22} & a_{23} \\ 0 & 0 & 0 & a_{33} \end{bmatrix} \begin{matrix} i \\ \vdots \\ n-i-1 \end{matrix}$$

$$U = [a_{00} \ a_{01} \ a_{11} \ a_{02} \ a_{12} \ a_{22} \ a_{03} \ a_{13} \ a_{23} \ a_{33}]$$

(a) (8 points) What is the number of items stored in U if A has n rows and n columns?

4

$$(N + (N-1) + \dots + 1) = \frac{N \cdot (N+1)}{2} = \binom{N+1}{2}$$

(b) (4 points) What is x if $U[x]$ stores the entry $A[i][j]$? Show your calculations.

5

Answer 1: $n + (n-1) + \dots + (n-(i-1)) + (j-i) = n \cdot i - \frac{(i-1) \cdot i}{2} + j - i$

Answer 2: $\binom{n+1}{2} - \binom{n-i+1}{2} - (n-j) = (n-1)i - \binom{i}{2} + j$

6

(c) (8 points) Please fill in the blanks in the method `readtriangularmatrix(int[], int)` that reads integers from the keyboard and fills the one-dimensional array U with entries of A as shown in the example.

```
void readtriangularmatrix(int U[], int n)
{
    int i, j, k;
    if(n*(n+1)\ 2 > MAX_SIZE){
        printf("\n invalid array size \ n");
        exit(-1);
    }
    else
        for(i=0; i<=n-1; i++){
            k = .....
            for(j=0; j<=i; j++)
                scanf("%d", ..... );
        }
}
```

$\rightarrow \binom{n+1}{2} - \binom{n-i+1}{2}; \text{ (I)}$

$\rightarrow \&U[k - (n-j)]; \text{ (II)}$

Or according to Answer 1:

(I): $(n-1)i - \binom{i}{2};$

(II): $\&u[k+j]$

10. (9 points) Assume a special type of matrix, named *stripe matrix*, is a matrix that has value 0 in all even-numbered columns (column 2, column 4, ...) as shown in the example below. (For simplicity, assume that all stripe matrices are N by N matrices and N is even.)

$$A = \begin{matrix} & \begin{matrix} j=0 & j=1 & j=2 & j=3 \end{matrix} \\ \begin{matrix} i=0 \\ i=1 \\ i=2 \\ i=3 \end{matrix} & \begin{bmatrix} 2 & 0 & 3 & 0 \\ 4 & 0 & 1 & 0 \\ 5 & 0 & -3 & 0 \\ 20 & 0 & 12 & 0 \end{bmatrix} \end{matrix}$$

- (a) (5 points) Write a 1-dimensional array representation of the matrix above, by filling in the entries of the array U .

$$U = \begin{array}{|c|c|c|c|c|c|c|c|c|c|} \hline & U[0] & U[1] & & & & & U[7] & & \\ \hline & 2 & 4 & 5 & 20 & 3 & 1 & -3 & 12 & \times & \times \\ \hline \end{array}$$

- (b) (4 points) What is the value of x if $U[x]$ stores a_{ij} ?

$$x = n \cdot \frac{j}{2} + i - 1$$

Check: a_{30} is at $U[3] = 20 \checkmark$

a_{32} is at $U[4+3] = U[7] = 12 \checkmark$

11. (8 points) The following operations are applied to an empty queue consecutively:
 Enqueue(7), Enqueue(5), Enqueue(2), Dequeue(), Enqueue(9), Enqueue(8),
 Enqueue(11), Enqueue(4), Enqueue(6), Enqueue(3), Dequeue(), Enqueue(4),
 Enqueue(3)

- (a) (4 points) Given an array `int A[10]`, fill the entries of the array representing the queue obtained above. (Assume that the array is initially empty.) Can all operations be performed?

A[0]	A[1]	A[2]	A[3]	A[4]	A[5]	A[6]	A[7]	A[8]	A[9]
*	5	2	9	8	11	4	6	3	4

(We cannot insert 3.)

- (b) (4 points) Given a *circular* array `int A[10]`, fill the entries of the circular array representing the queue obtained above. (Assume that the array is initially empty.) Can all operations be performed?

