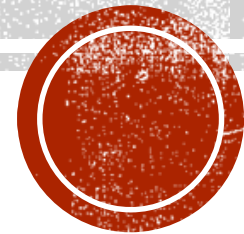


STACKS & QUEUES



STACKS & QUEUES

Two abstract data types

- Can be implemented using arrays or linked lists
- Restricted lists (only a small number of operations are performed on them)

Stacks

“Last In First Out” (LIFO)

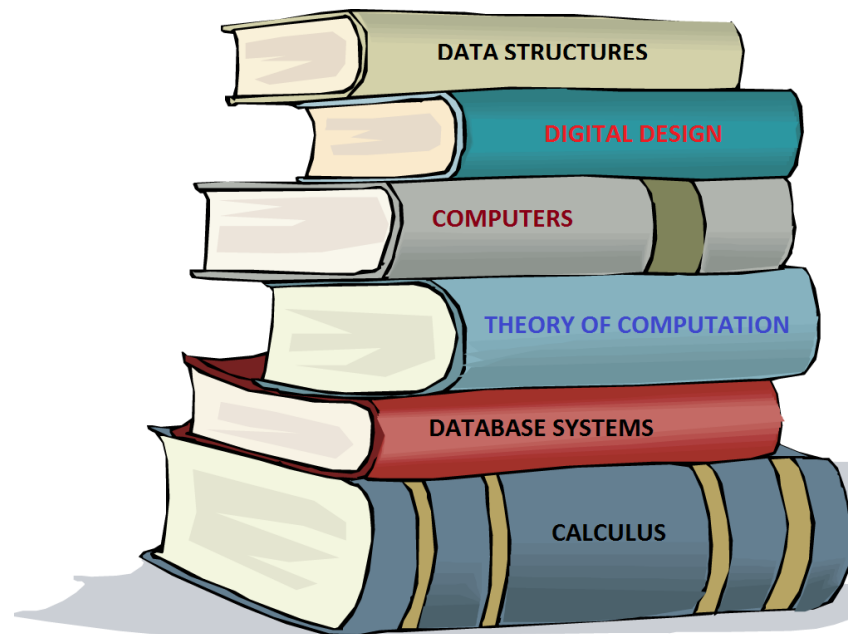
Queues

“First In First Out” (FIFO)



STACKS

A stack is an ordered lists in which insertions and deletions are made at one end called the *top*.



OPERATIONS ON STACKS

PUSH

Place an item on the stack

Peek

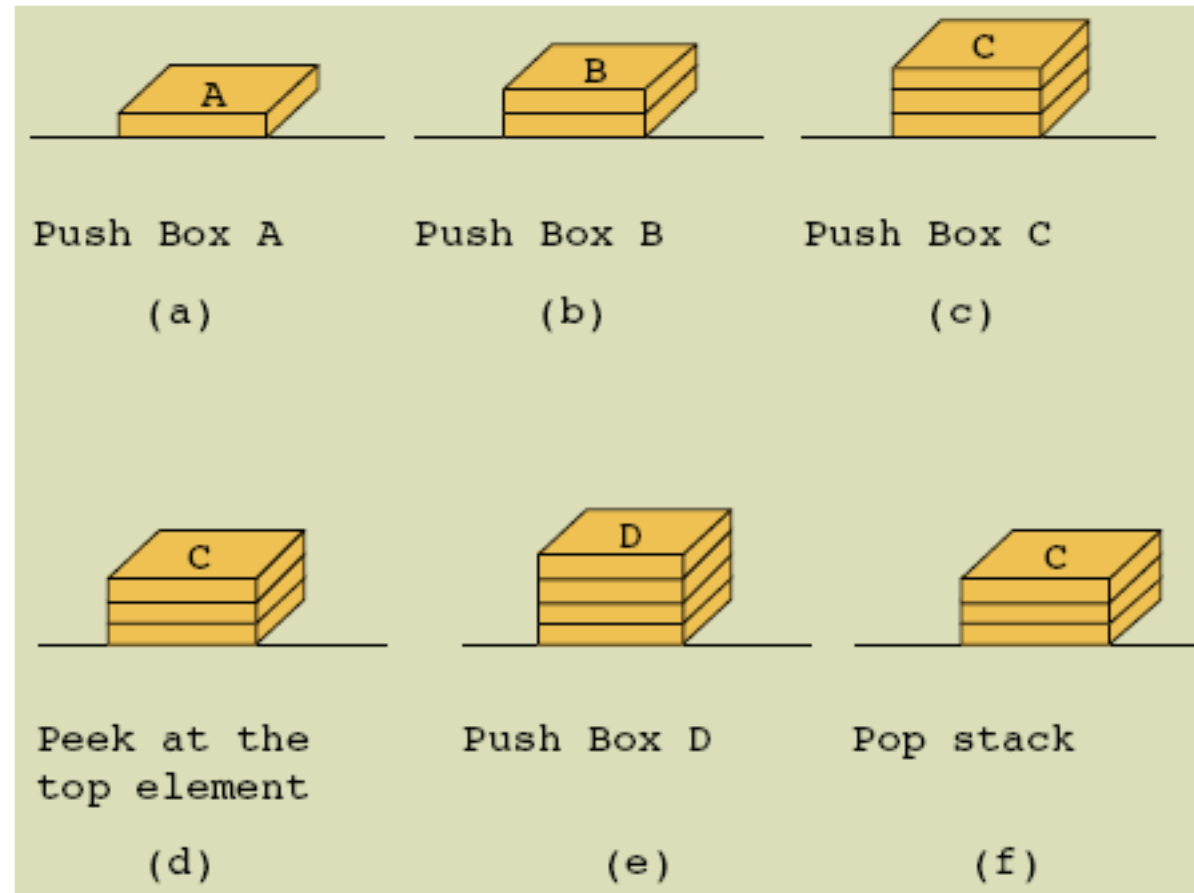
Look at the item on top of the stack, but do not remove it

POP

Look at the item on top of the stack and remove it



OPERATIONS ON STACKS



SOME APPLICATIONS

Converting a decimal number into a binary

Towers of Hanoi

Program execution

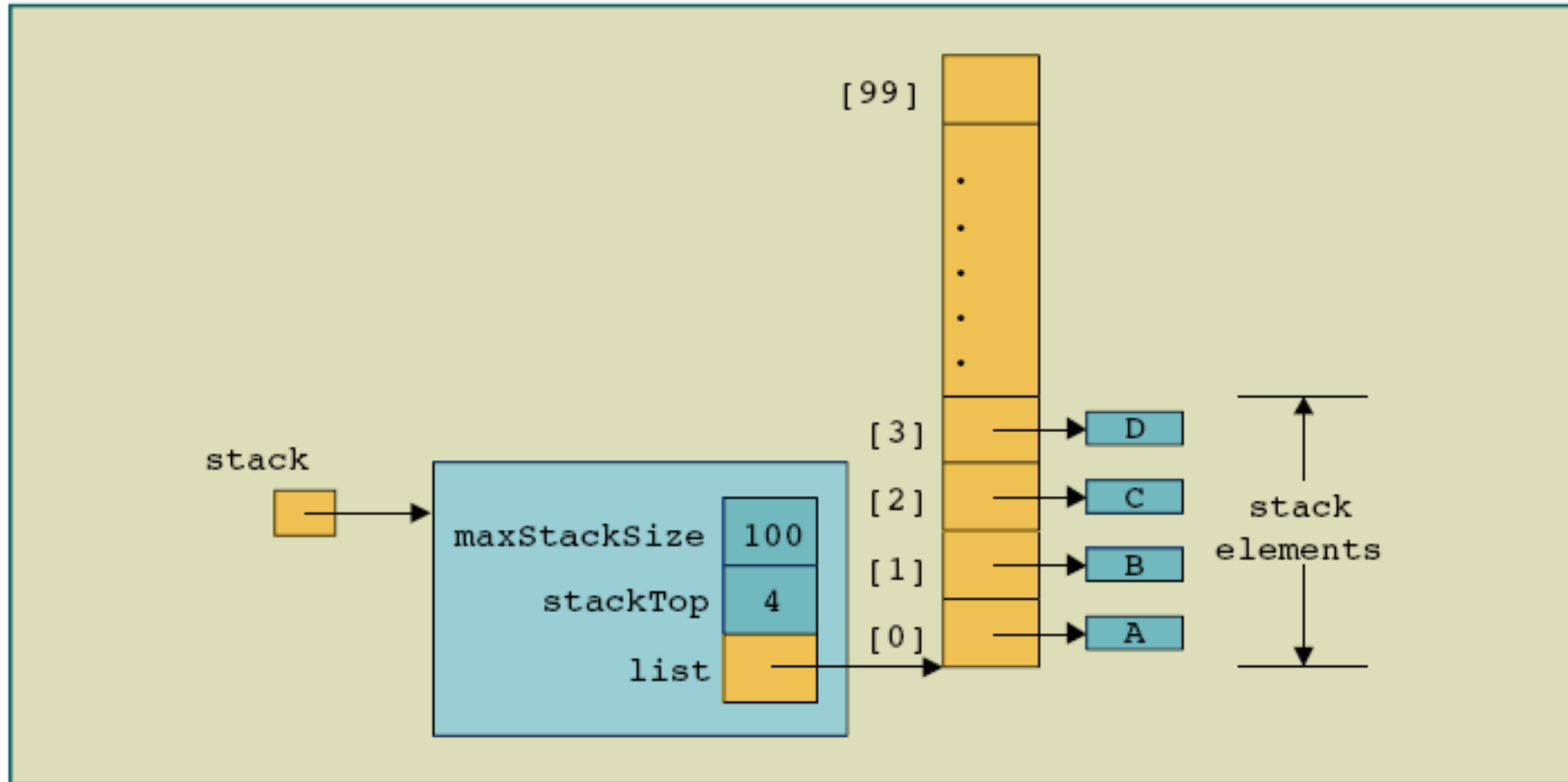
Parsing

Evaluating postfix expressions

....



IMPLEMENTATION OF STACKS USING ARRAYS



STACK

```
#define MAX_STACK_SIZE 100

/* stack structure*/
struct stack {
    int list[MAX_STACK_SIZE];
    int stackTop;
} st;
```



INITIALIZE STACK

```
void initializeStack()
{
    int i;
    for (i = 0 ; i < MAX_STACK_SIZE;
        st.list[i] = 0;
    st.stackTop = 0;
}
```



EMPTY STACK

```
void isEmpty()  
{  
    if( st.stackTop == 0)  
        return 1;  
    else  
        return 0;  
}
```



FULL STACK

```
void isFull()
{
    if( st.stackTop >= MAX_STACK_SIZE)
        return 1;
    else
        return 0;
}
```



PUSH STACK

```
void push(int item)
{
    if (!isFull()){
        st.list[st.stackTop] = item;
        st.stackTop++;
    }
    else
        printf("Stack Is Full!");
}
```



PEEK STACK

```
int peek()
{
    if (!isEmpty()){
        return st.list[st.stackTop-1] ;
    }
    else
        return -1;
}
```



POP STACK

```
void pop()
{
    if (!isEmpty()){
        st.stackTop--;
        st.list[st.stackTop] = 0;
    }
    else
        printf("Stack Is Empty!");
}
```



CONVERTING A DECIMAL NUMBER TO A BINARY

```
function outputInBinary(Integer n)
    Stack s = new Stack
    while n > 0 do
        Integer bit = n modulo 2
        s.push(bit)
        if s is full
            then return error
        end if
        n = floor(n / 2)
    end while
    while s is not empty do
        output(s.pop())
    end while
end function
```

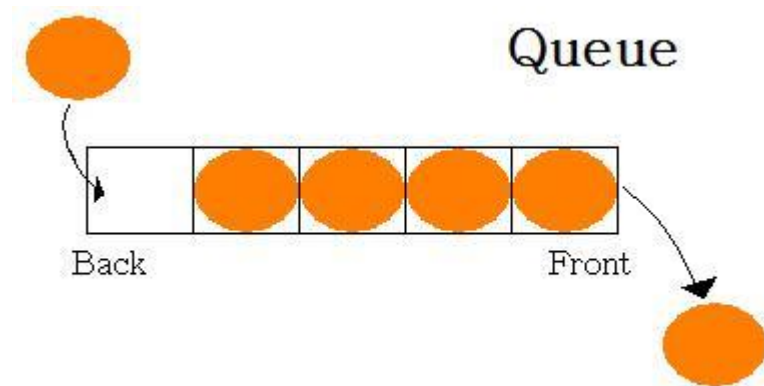


QUEUES

A queue is an ordered list in which all insertions take place at one end called the *rear/back* and all deletions take place at the opposite end called the *front*.

A queue is a First In First Out data structure.

As in a stack, the middle elements of the queue are inaccessible.



CONCEPTUAL QUEUE



QUEUE or STACK ?



QUEUE

```
#define MAX_QUEUE_SIZE 100
/* queue structure*/
struct stack {
    int key;
} element;
element queue[MAX_QUEUE_SIZE];
int rear = -1;
int front = -1;
```



INITIALIZE QUEUE

```
void initializeQueue()
{
    int i;
    for (i = 0 ; i < MAX_QUEUE_SIZE;
        i++)
        queue[i].key = 0;
    rear = front = -1;
}
```



EMPTY QUEUE

```
void isEmpty()  
{  
    if( front == rear)  
        return 1;  
    else  
        return 0;  
}
```



FULL QUEUE

```
void isFull()
{
    if( rear == MAX_QUEUE_SIZE)
        return 1;
    else
        return 0;
}
```



ADD QUEUE

```
void addq(int *read, element item)
{
    if(*rear == MAX_QUEUE_SIZE-1)
        then return error
    queue[++*read] = item;
}
```



DELETE QUEUE

```
int deleteq(int *front, int rear)
{
    if (*front == rear){
        return empty error;
    }
    return queue[++*front];
}
```



JOB SCHEDULING

<i>front</i>	<i>rear</i>	$Q[0]$	$Q[1]$	$Q[2]$	$Q[3]$	Comments
-1	-1					queue is empty
-1	0	J1				Job 1 is added
-1	1	J1	J2			Job 2 is added
-1	2	J1	J2	J3		Job 3 is added
0	2		J2	J3		Job 1 is deleted
1	2			J3		Job 2 is deleted



CIRCULAR QUEUE

When the queue is full (the rear index is equals `MAX_QUEUE_SIZE`)
we should move the entire queue to the left and
recalculate the rear

→ The first element of the queue is at `queue[0]` and front is at `-1`.

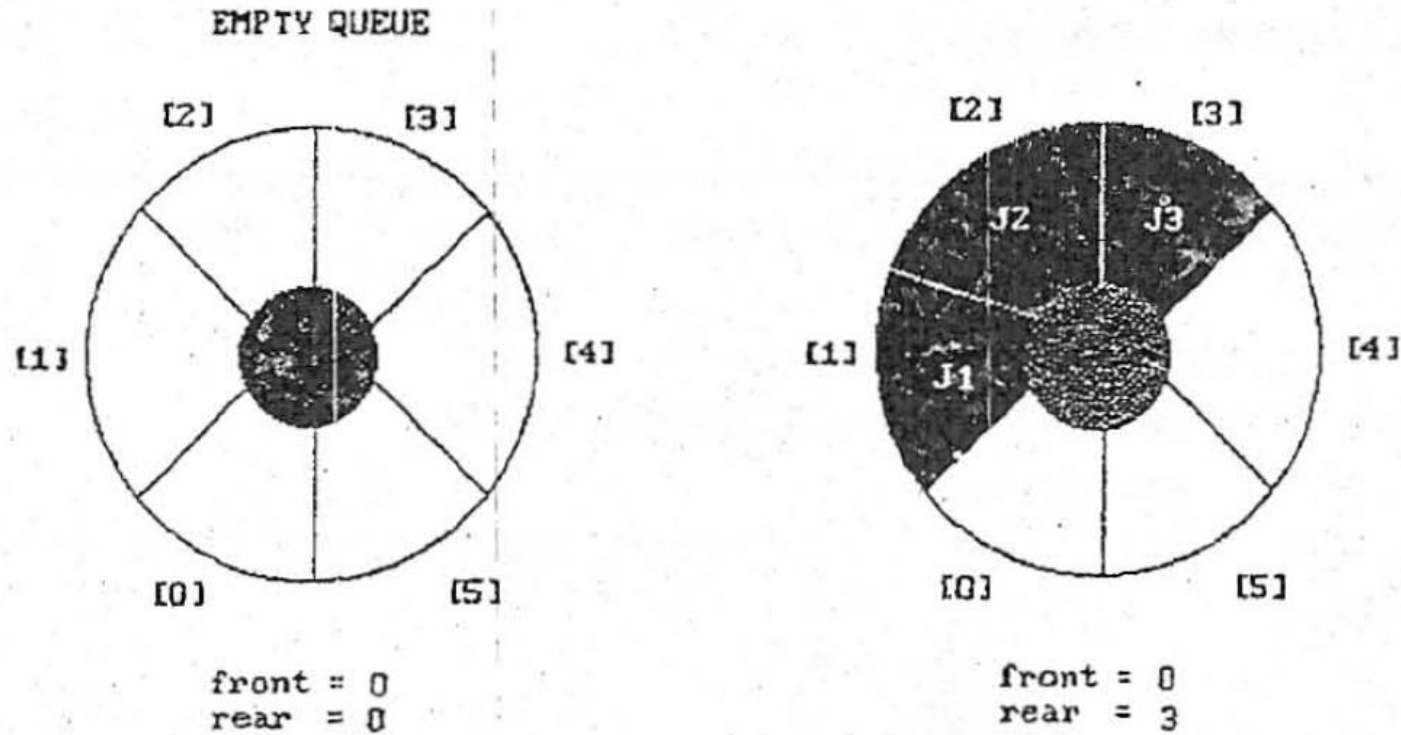
Shifting an array is time-consuming

$O(\text{MAX_QUEUE_SIZE})$

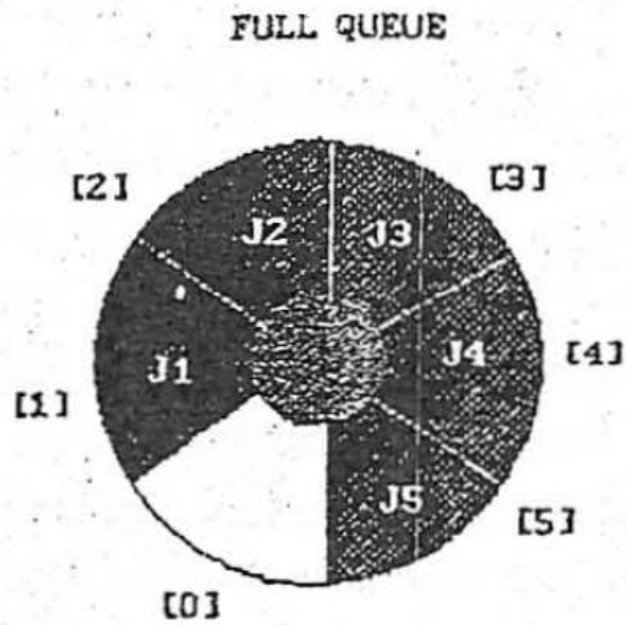


CIRCULAR QUEUE

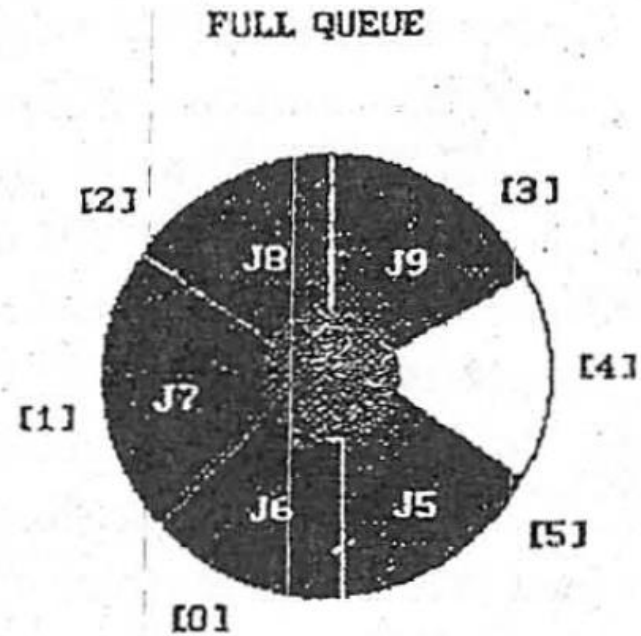
More efficient queue representation



FULL CIRCULAR QUEUE



front = 0
rear = 5



front = 4
rear = 3



INITIALIZE CIRCULAR QUEUE

```
void initializeQueue()
{
    int i;
    for (i = 0 ; i < MAX_QUEUE_SIZE;
        i++)
        queue[i].key = 0;
    rear = front = 0;
}
```



EMPTY QUEUE

```
void isEmpty()  
{  
    if( front == rear)  
        return 1;  
    else  
        return 0;  
}
```



ADD CIRCULAR QUEUE

```
void addq(int front, int *rear, element item)
{
    *rear = (*rear + 1) % MAX_QUEUE_SIZE
    if(front == *rear){
        queueIsFull(rear);
        return;
    }
    queue[*rear] = item;
}
```

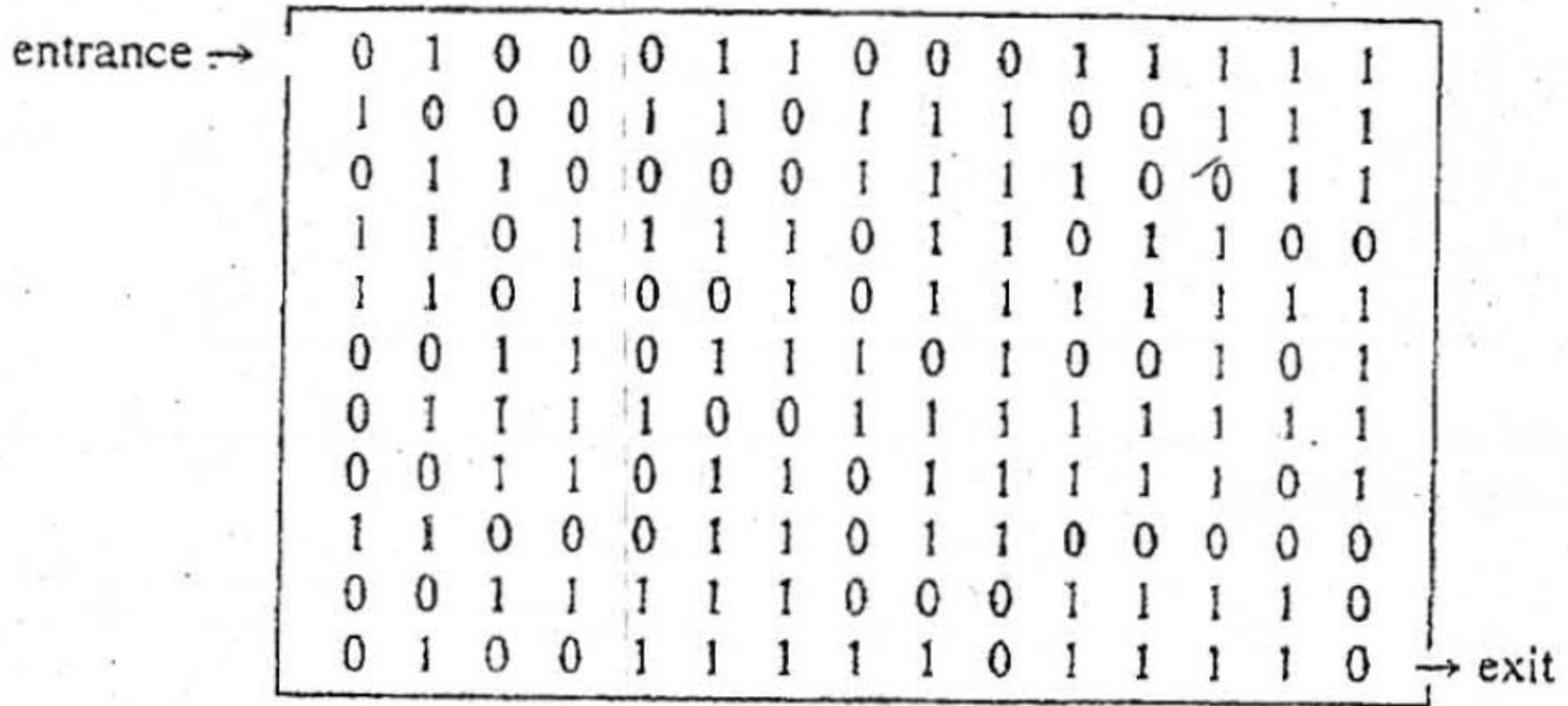


DELETE CIRCULAR QUEUE

```
int deleteq(int *front, int rear)
{
    if (*front == rear){
        return empty error;
        *front = (*front + 1) %
MAX_QUEUE_SIZE
        return queue[*front];
    }
```



A MAZING PROBLEM

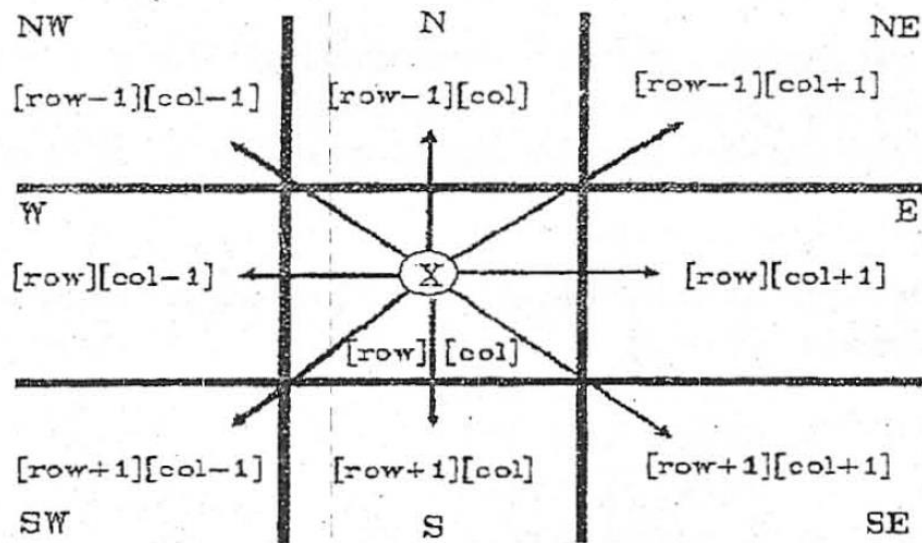


DIRECTIONS

```
typedef struct {  
    short int vert;  
    short int horiz;  
} offsets;  
offsets move[8];
```



ALLOWABLE MOVES



Name	Dir	<i>move[dir].vert</i>	<i>move[dir].horiz</i>
N	0	-1	0
NE	1	-1	1
E	2	0	1
SE	3	1	1
S	4	1	0
SW	5	1	-1
W	6	0	-1
NW	7	-1	-1

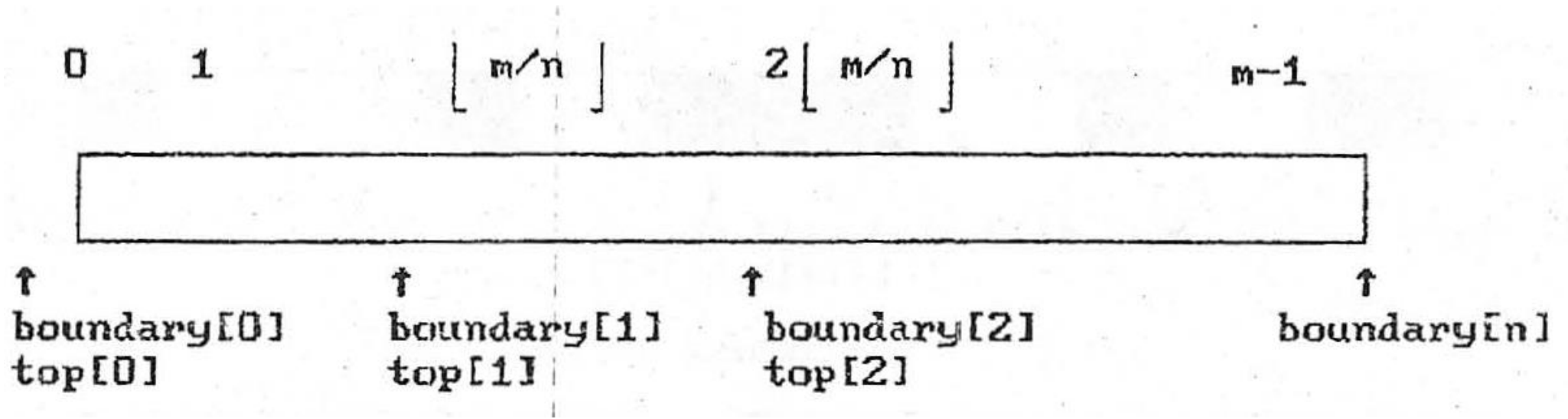
`next_row = row + move[dir].vert;`
`next_col = col + move[dir].horiz;`



IMPLEMENTATION



MULTIPLE STACKS & QUEUES



MULTIPLE STACKS

```
#define MEMORY_SIZE 100
#define MAX_STACKS 10
element memory[MEMORY_SIZE];
int top[MAX_STACKS];
int boundary[MAX_STACKS];
int n;          //number of stacks entered by the user

top[0] = boundary[0] = -1;
for(i = 1; i < n ; i++)
    top[i] = boundary[i] = (MEMORY_SIZE/n)*i;
boundary[n] = MEMORY_SIZE - 1;
```



MULTIPLE STACKS - ADD

```
void add(int i, element item)
{
    if(top[i] == boundary[i+1])
        stack_full(i);
    memory[++top[i]] = item;
}
```



MULTIPLE STACKS - DELETE

```
element delete(int i)
{
    if(top[i] == boundary[i])
        return stack_empty(i);
    return memory[top[i]--];
}
```



MULTIPLE STACKS

1. $\text{stack_no} < j < n$, such that there is free space between stacks j and $j+1$

move stacks $\text{stack_no}+1, \text{stack_no}+2, \dots, j$ one position right.
this creates a space between stack_no and $\text{stack_no}+1$.

2. $0 \leq j < \text{stack_no}$, such that there is free space between stacks j and $j+1$

move stacks $j+1, j+2, \dots, \text{stack_no}$ one position right.
this also creates a space between stack_no and $\text{stack_no}+1$.

3. No j satisfying either condition 1 or condition 2, then there is no free space.

$O(\text{MEMORY_SIZE})$

